

# Graphs and Genomes

Michael Schatz

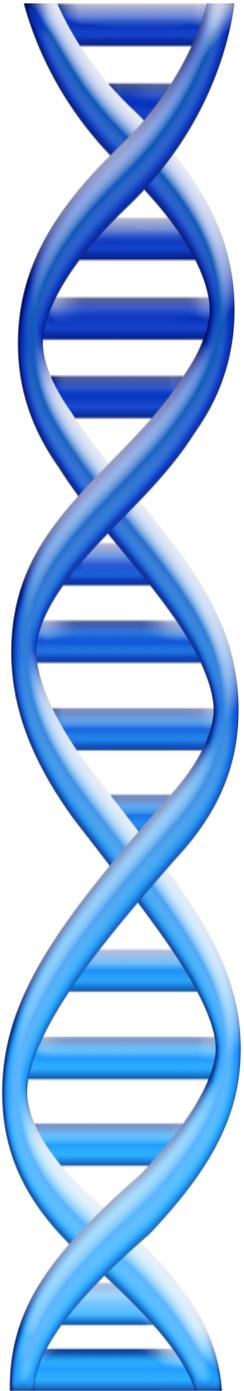
July 27, 2012

CSHL Undergraduate Research Program



# Outline

1. Graph Searching
2. Assembly by analogy
3. Genome Assembly



# Biological Networks

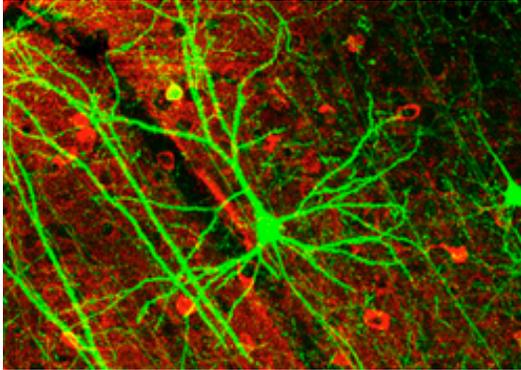
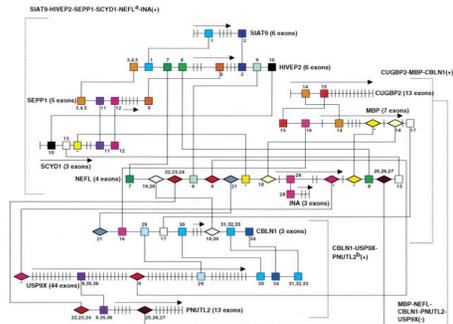
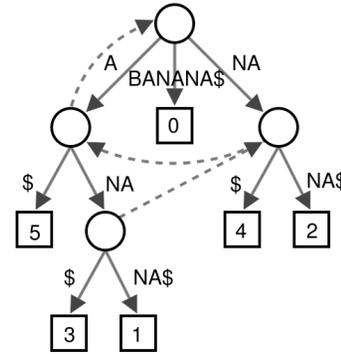
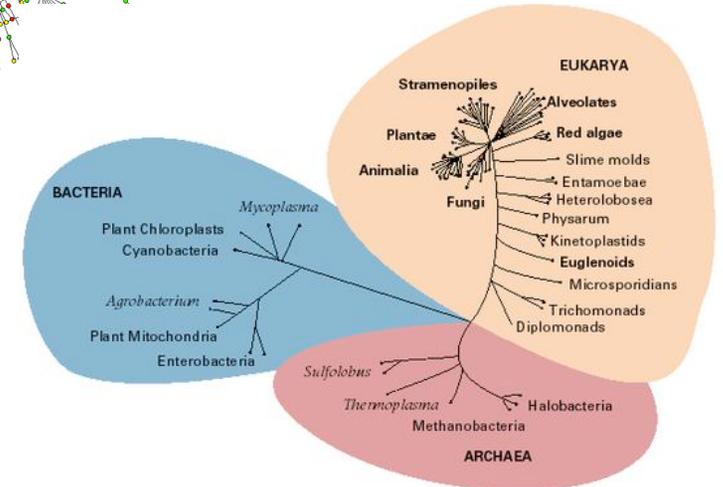
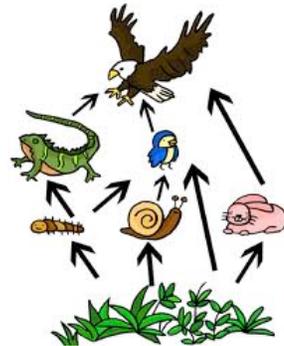
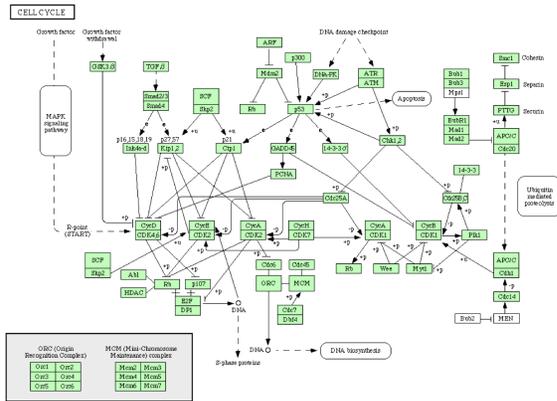
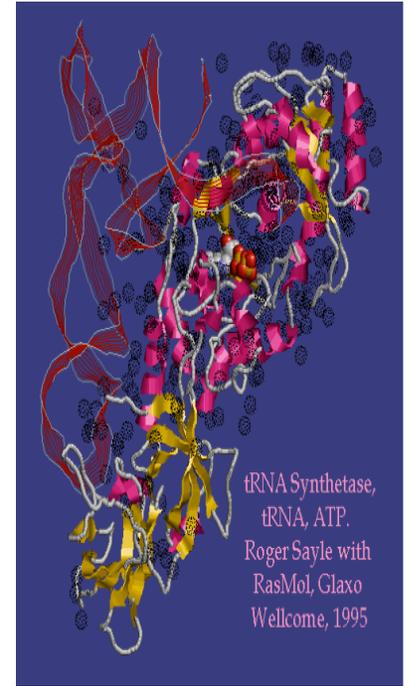
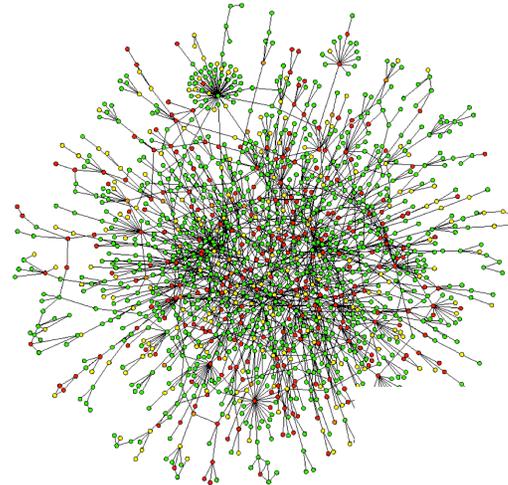


Figure 5 Putative regulatory elements shared between groups of correlated and anticorrelated genes

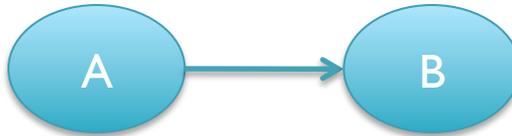


Vanessa M. Brown et al. Genome Res. 2002; 12: 868-884

Cold Spring Harbor Laboratory Press



# Graphs

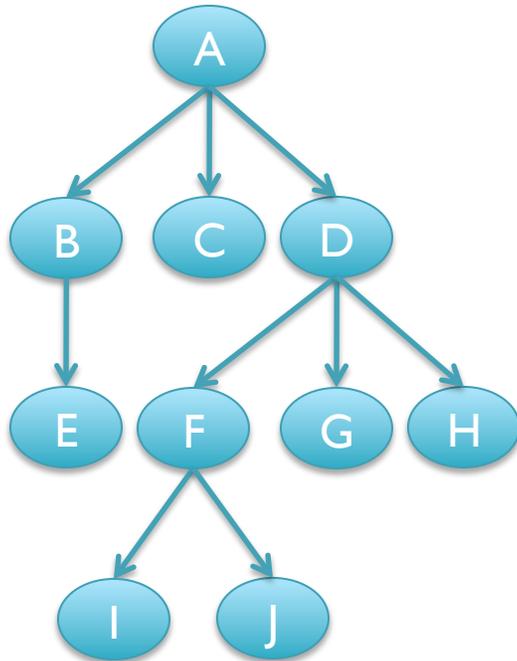


- Nodes
  - People, Proteins, Genes, Neurons, Sequences, Numbers, ...
- Edges
  - A is connected to B
  - A is related to B
  - A regulates B
  - A precedes B
  - A interacts with B
  - A activates B
  - ...

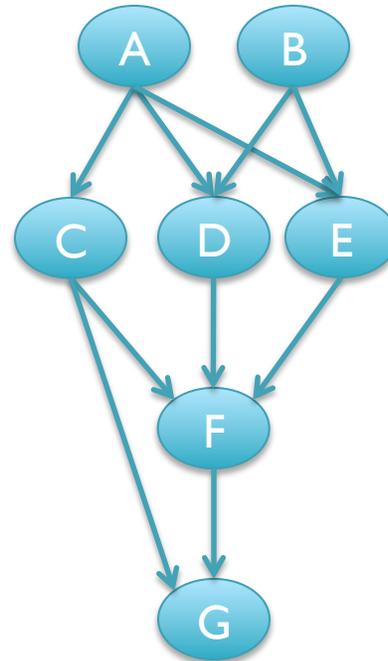
# Graph Types



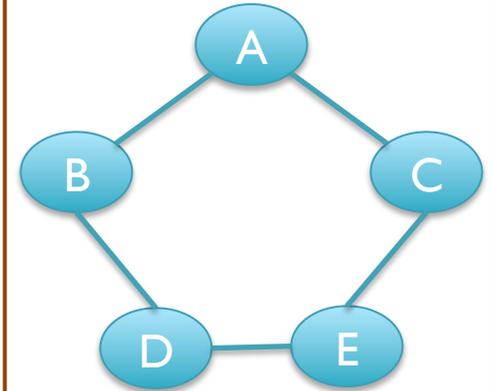
List



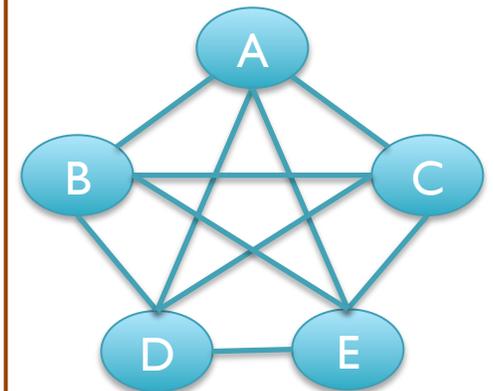
Tree



Directed  
Acyclic  
Graph



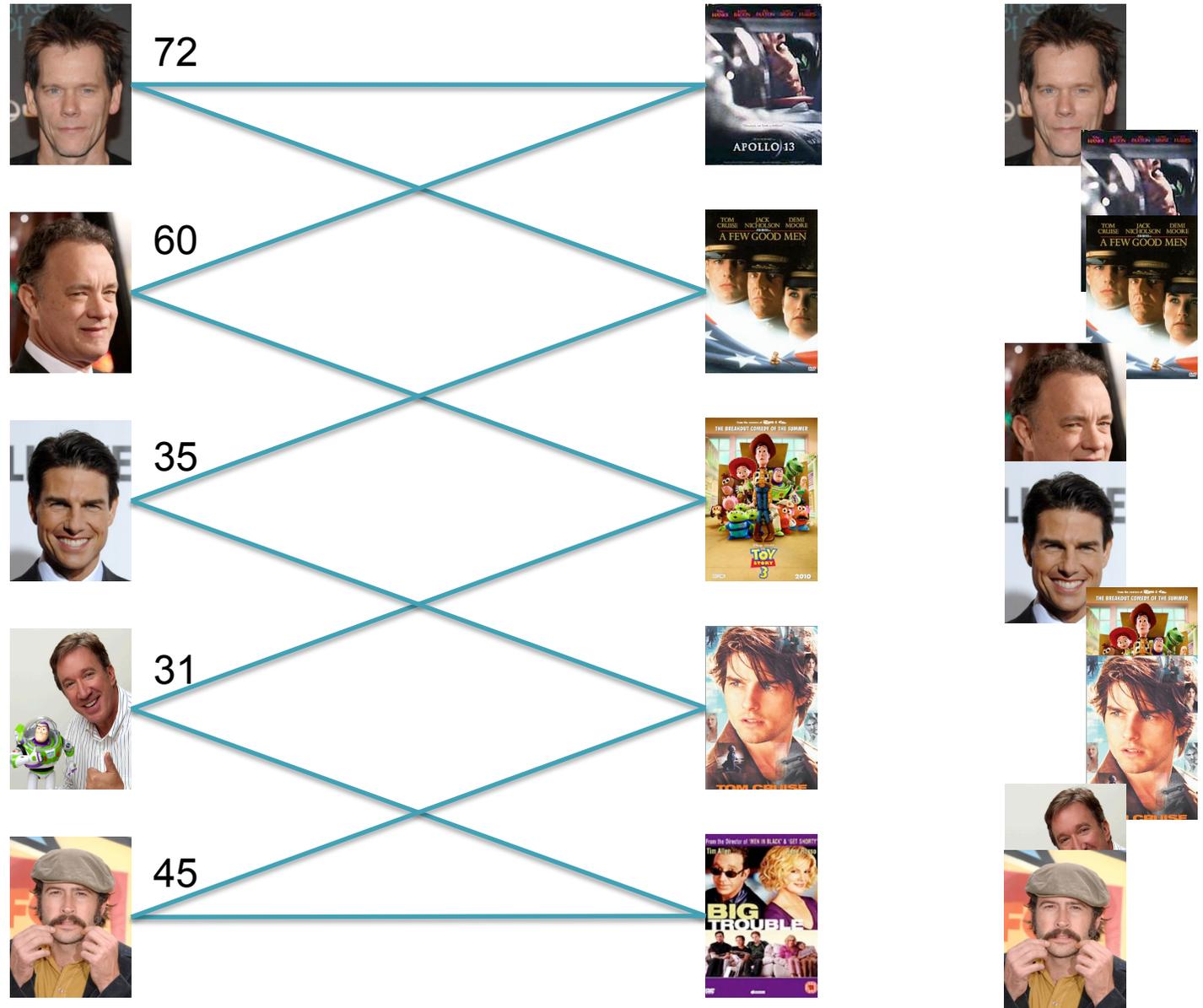
Cycle



Complete

# Kevin Bacon and Bipartite Graphs

Find the **shortest** path from Kevin Bacon to Jason Lee



Breadth First Search:  
4 hops

Bacon Distance:  
2

# BFS and TSP

- BFS computes the shortest path between a pair of nodes in  $O(|E|) = O(|N|^2)$
- What if we wanted to compute the shortest path visiting every node once?
  - Traveling Salesman Problem

$$\text{ABDCA: } 4+2+5+3 = 14$$

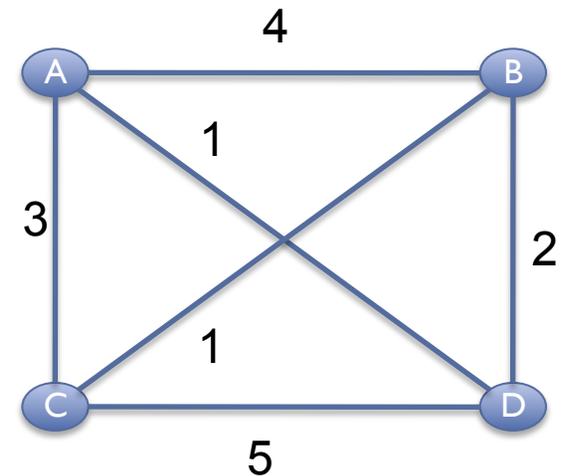
$$\text{ACDBA: } 3+5+2+4 = 14^*$$

$$\text{ABCD A: } 4+1+5+1 = 11$$

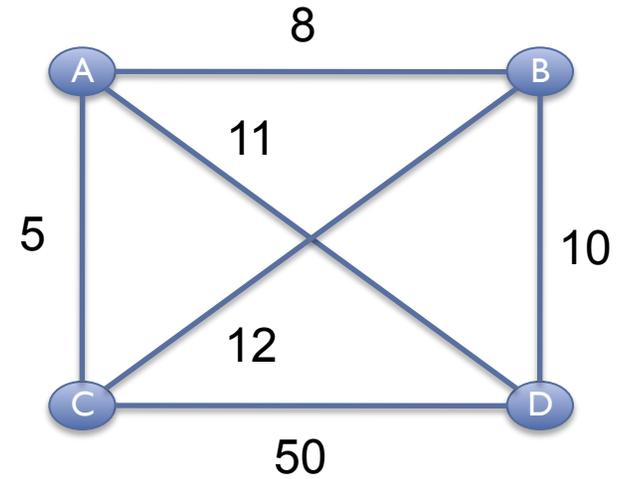
$$\text{ADCBA: } 1+5+1+4 = 11^*$$

$$\text{ACBDA: } 3+1+2+1 = 7$$

$$\text{ADBCA: } 1+2+1+3 = 7^*$$



# Greedy Search



# Greedy Search

## **Greedy Search**

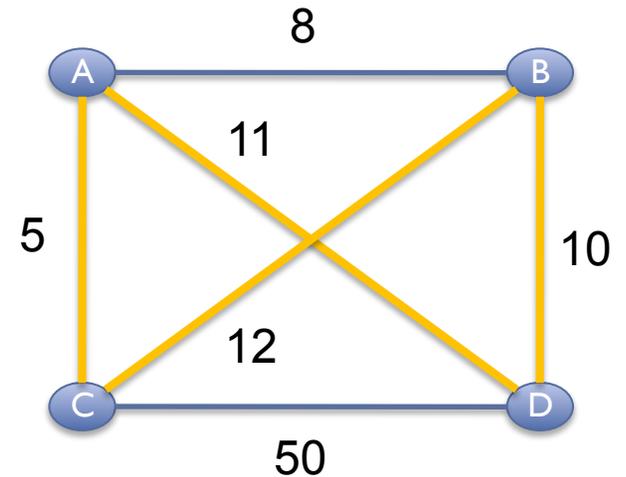
```
cur=graph.randNode()
```

```
while (!done)
```

```
    next=cur.getNextClosest()
```

Greedy:  $ABDCA = 5+8+10+50= 73$

Optimal:  $ACBDA = 5+11+10+12 = 38$



Greedy finds the global optimum only when

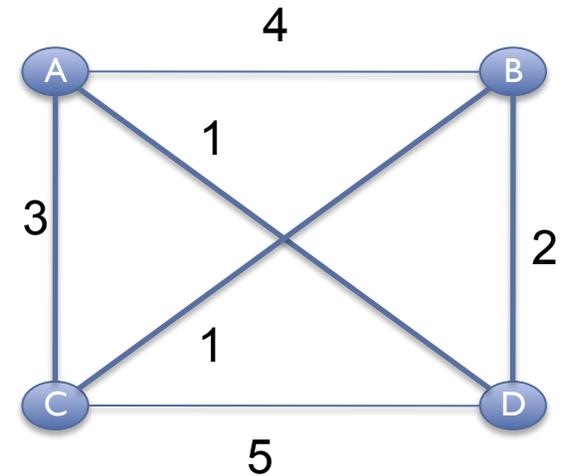
1. Greedy Choice: Local is correct without reconsideration
2. Optimal Substructure: Problem can be split into subproblems

Optimal Greedy: Making change with the fewest number of coins

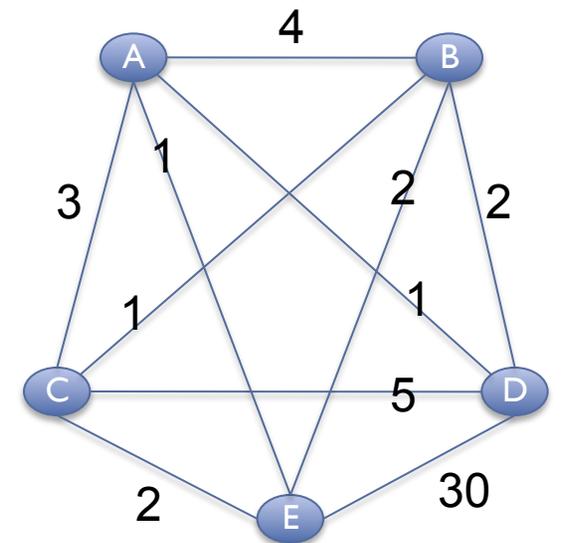
# TSP Complexity

- No fast solution
  - Knowing optimal tour through  $n$  cities doesn't seem to help much for  $n+1$  cities

[How many possible tours for  $n$  cities?]

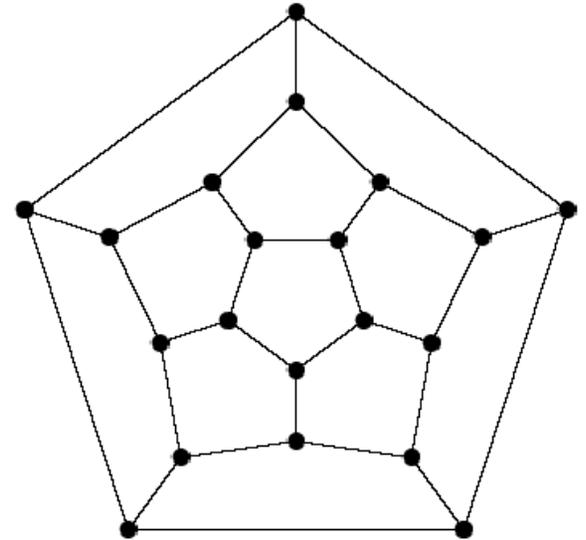


- Extensive searching is the only provably correct algorithm
  - Brute Force:  $O(n!)$ 
    - ~20 cities max
    - $20! = 2.4 \times 10^{18}$
  - Branch-and-Bound can sometimes help



# TSP and NP-complete

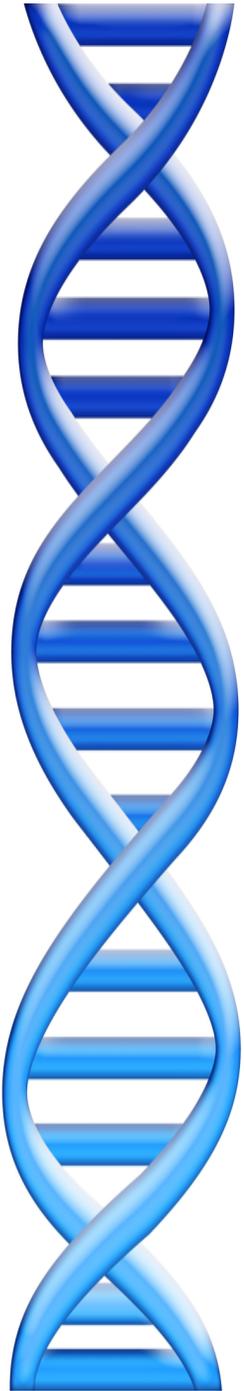
- TSP is one of many extremely hard problems of the class NP-complete
  - Extensive searching is the only way to find an exact solution
  - Often have to settle for approx. solution



- **WARNING:** Many biological problems are in this class
  - Find a tour that visits every node once (Genome Assembly)
  - Find the smallest set of vertices covering the edges (Essential Genes)
  - Find the largest clique in the graph (Protein Complexes)
  - Find the highest mutual information encoding scheme (Neurobiology)
  - Find the best set of moves in tetris
  - ...
  - [http://en.wikipedia.org/wiki/List\\_of\\_NP-complete\\_problems](http://en.wikipedia.org/wiki/List_of_NP-complete_problems)

# Outline

1. Graph Searching
2. **Assembly by analogy**
3. Genome Assembly



# Shredded Book Reconstruction

- Dickens accidentally shreds the first printing of A Tale of Two Cities
  - Text printed on 5 long spools

|        |     |      |      |        |        |     |     |       |       |        |        |     |     |     |     |         |         |     |     |     |     |                  |                  |
|--------|-----|------|------|--------|--------|-----|-----|-------|-------|--------|--------|-----|-----|-----|-----|---------|---------|-----|-----|-----|-----|------------------|------------------|
| It was | the | best | of   | times, | it     | was | the | worst | of    | times, | it     | was | the | age | of  | wisdom, | it      | was | the | age | of  | foolishness, ... |                  |
| It was | the | best | of   | times, | it     | was | the | worst | of    | times, | it     | was | the | age | of  | wisdom, | it      | was | the | age | of  | foolishness, ... |                  |
| It was | the | best | of   | times, | it     | was | the | worst | of    | times, | it     | was | the | age | of  | wisdom, | it      | was | the | age | of  | foolishness, ... |                  |
| It was | the | best | of   | times, | it     | was | the | worst | of    | times, | it     | was | the | age | of  | wisdom, | it      | was | the | age | of  | foolishness, ... |                  |
| It     | was | the  | best | of     | times, | it  | was | the   | worst | of     | times, | it  | was | the | age | of      | wisdom, | it  | was | the | age | of               | foolishness, ... |

- How can he reconstruct the text?
  - 5 copies x 138,656 words / 5 words per fragment = 138k fragments
  - The short fragments from every copy are mixed together
  - Some fragments are identical

# Greedy Reconstruction

It was the best of  
age of wisdom, it was  
best of times, it was  
it was the age of  
it was the age of  
it was the worst of  
of times, it was the  
of times, it was the  
of wisdom, it was the  
the age of wisdom, it  
the best of times, it  
the worst of times, it  
times, it was the age  
times, it was the worst  
was the age of wisdom,  
was the age of foolishness,  
was the best of times,  
was the worst of times,  
wisdom, it was the age  
worst of times, it was

It was the best of  
was the best of times,  
the best of times, it  
best of times, it was  
of times, it was the  
of times, it was the  
times, it was the worst  
times, it was the age

The repeated sequence make the correct reconstruction ambiguous

- It was the best of times, it was the [worst/age]

Model sequence reconstruction as a graph problem.

# de Bruijn Graph Construction

- $D_k = (V, E)$ 
  - $V =$  All length- $k$  subfragments ( $k < l$ )
  - $E =$  Directed edges between consecutive subfragments
    - Nodes overlap by  $k-1$  words

Original Fragment

It was the best of

Directed Edge

It was the best → was the best of

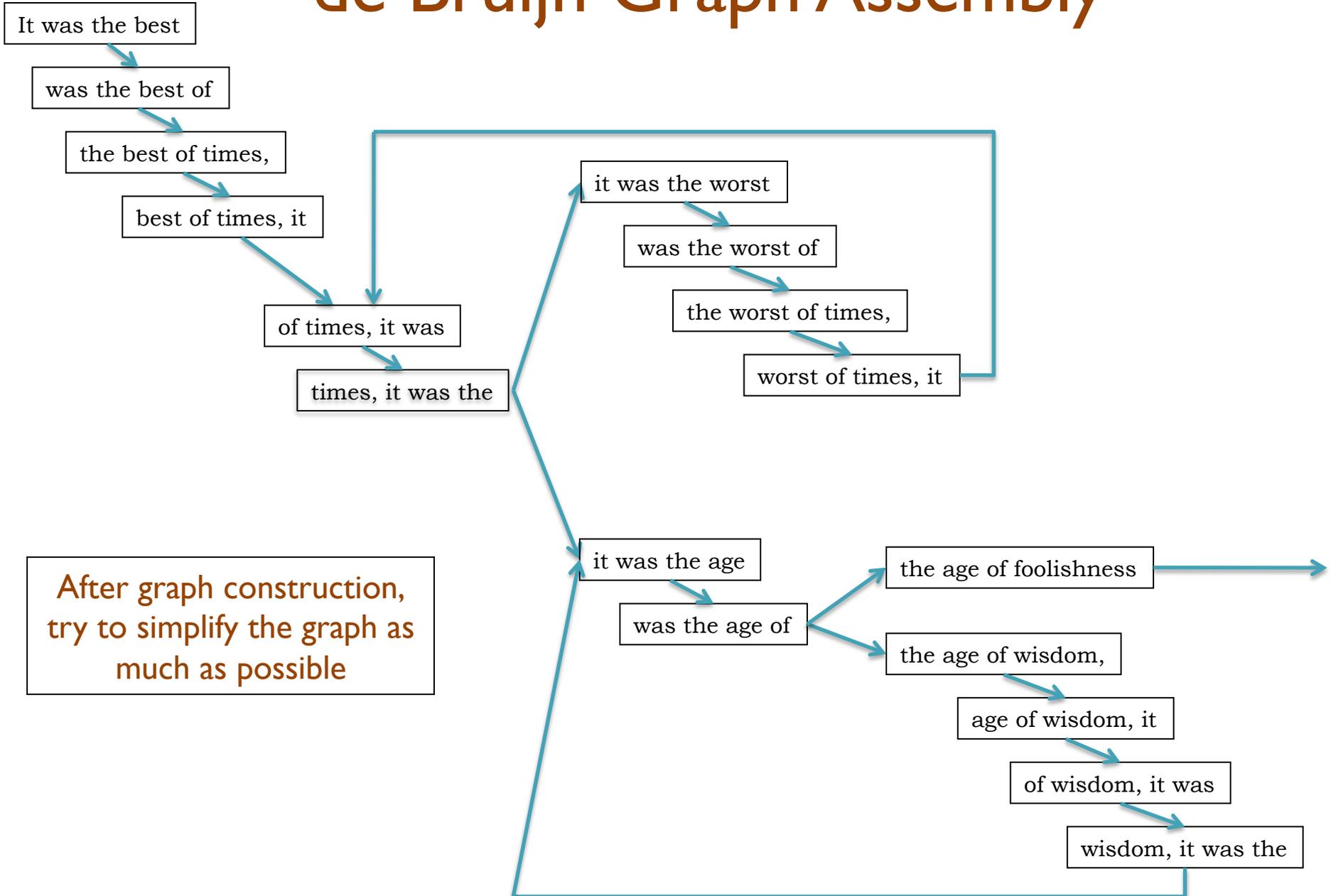
- Locally constructed graph reveals the global sequence structure
  - Overlaps between sequences implicitly computed

de Bruijn, 1946

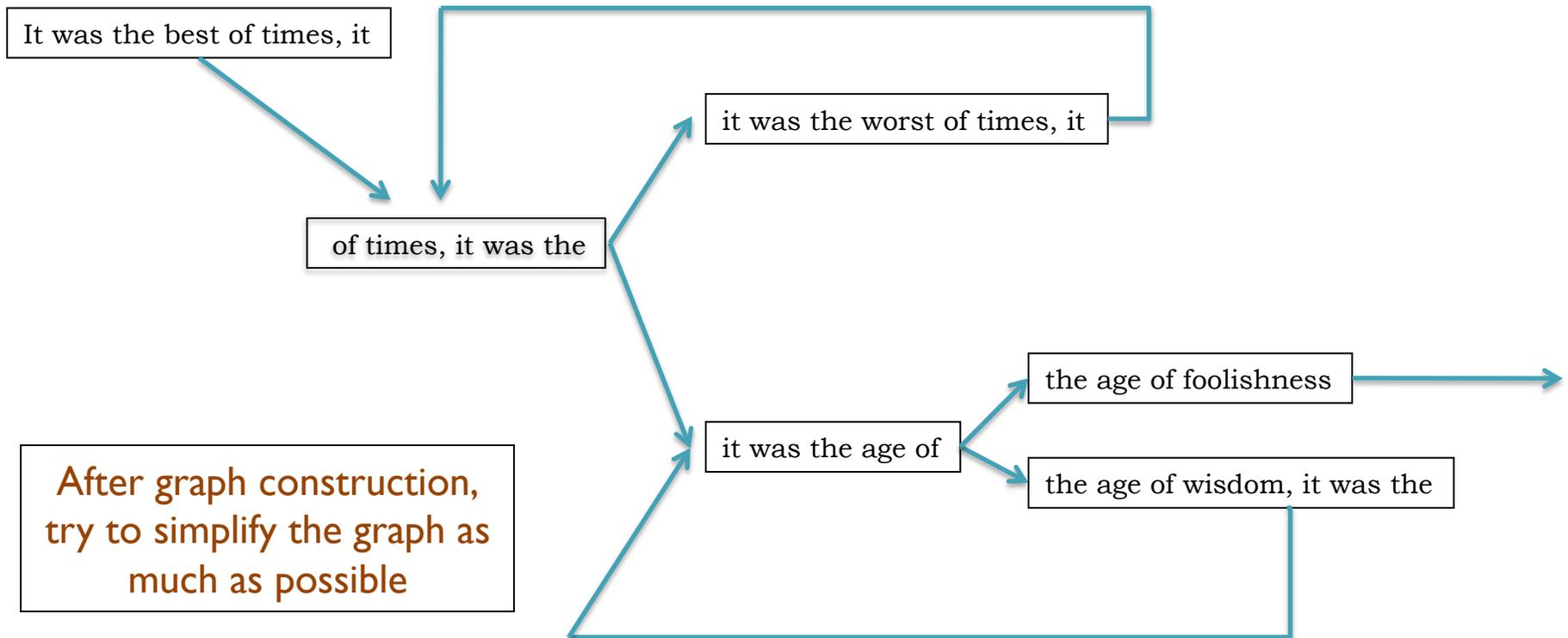
Idury and Waterman, 1995

Pevzner, Tang, Waterman, 2001

# de Bruijn Graph Assembly

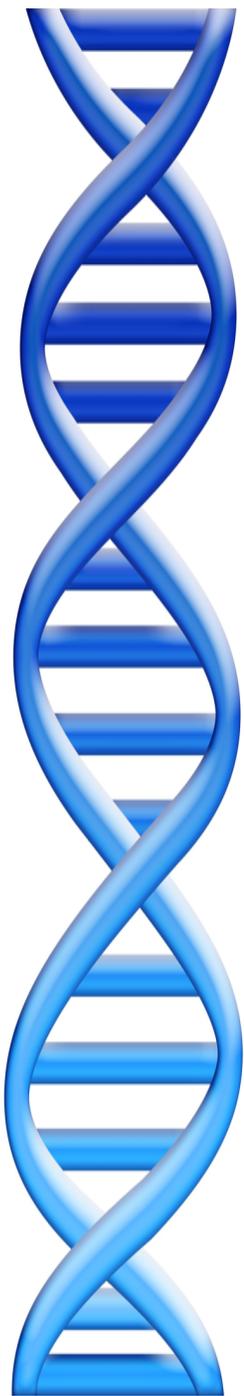


# de Bruijn Graph Assembly

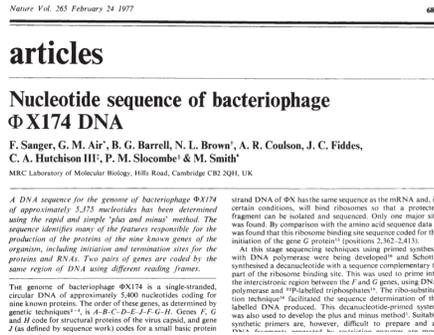


# Outline

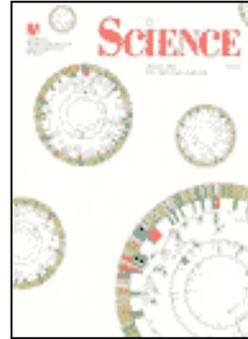
1. Genome Assembly by Analogy
2. Graph Searching
3. **Genome Assembly**



# Milestones in Genome Assembly



1977. Sanger *et al.*  
1<sup>st</sup> Complete Organism  
5375 bp



1995. Fleischmann *et al.*  
1<sup>st</sup> Free Living Organism  
TIGR Assembler. 1.8Mbp



1998. C. elegans SC  
1<sup>st</sup> Multicellular Organism  
BAC-by-BAC Phrap. 97Mbp



2000. Myers *et al.*  
1<sup>st</sup> Large WGS Assembly.  
Celera Assembler. 116 Mbp



2001. Venter *et al.*, IHGSC  
Human Genome  
Celera Assembler/GigaAssembler. 2.9 Gbp



2010. Li *et al.*  
1<sup>st</sup> Large SGS Assembly.  
SOAPdenovo 2.2 Gbp



Like Dickens, we must computationally reconstruct a genome from short fragments

# Current Applications

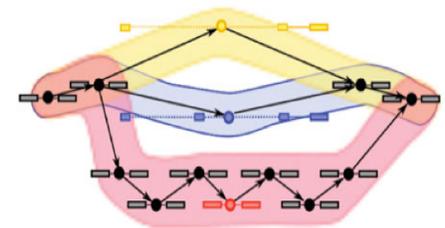
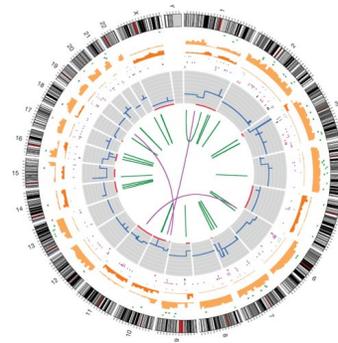
- Novel genomes



- Metagenomes

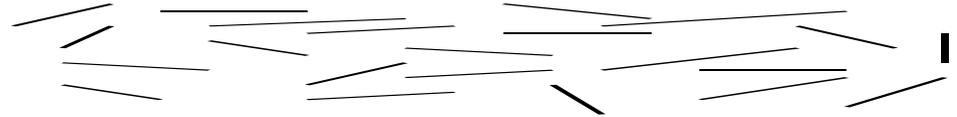


- Sequencing assays
  - Structural variations
  - Transcript assembly
  - ...



# Assembling a Genome

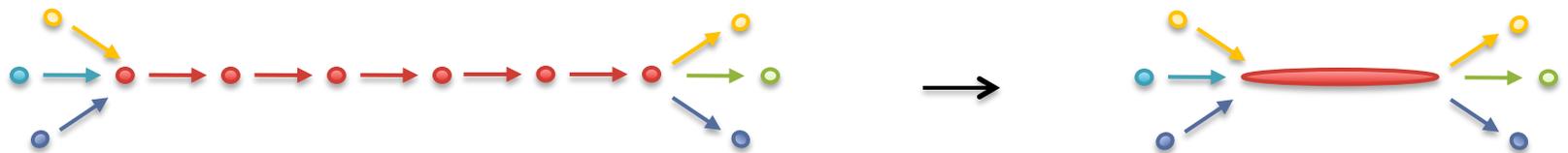
1. Shear & Sequence DNA



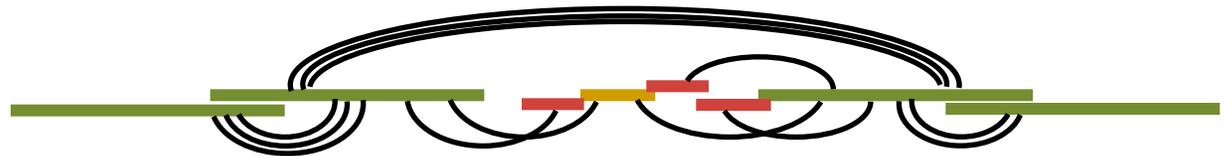
2. Construct assembly graph from overlapping reads

...AGCCTAGACCTACAGGATGCGCGACACGT  
GGATGCGCGACACGTTCGCATATCCGGT...

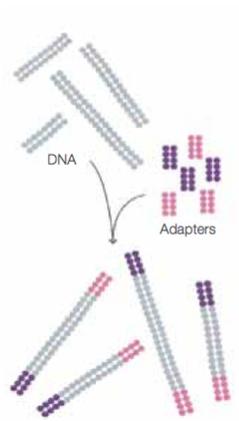
3. Simplify assembly graph



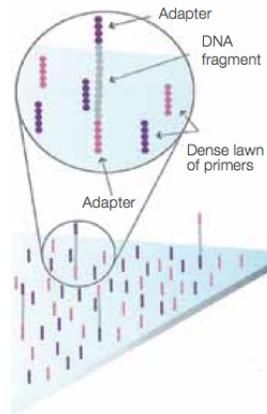
4. Detangle graph with long reads, mates, and other links



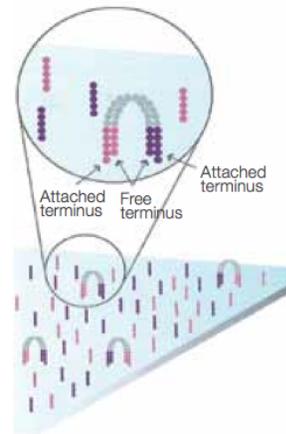
# Illumina Sequencing by Synthesis



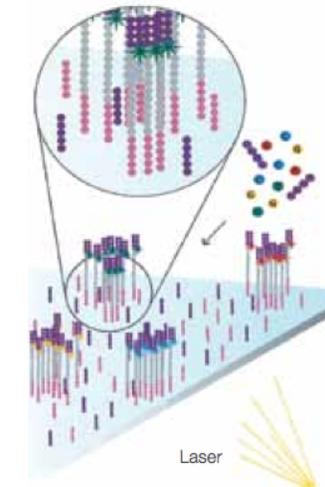
1. Prepare



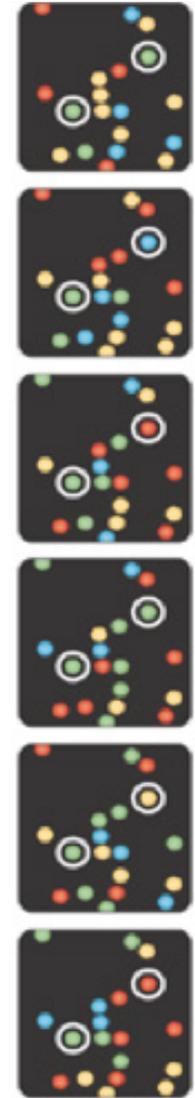
2. Attach



3. Amplify



4. Image



5. Basecall

Metzker (2010) Nature Reviews Genetics 11:31-46

[http://www.illumina.com/documents/products/techspotlights/techspotlight\\_sequencing.pdf](http://www.illumina.com/documents/products/techspotlights/techspotlight_sequencing.pdf)

# Paired-end and Mate-pairs

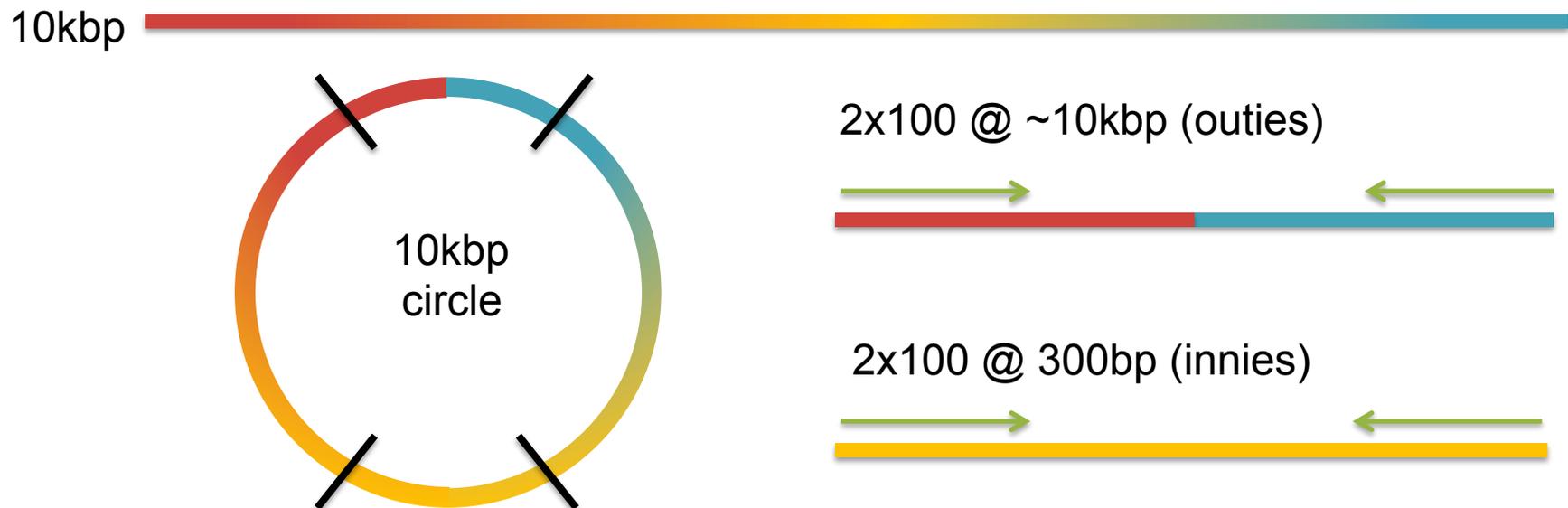
## ***Paired-end sequencing***

- Read one end of the molecule, flip, and read the other end
- Generate pair of reads separated by up to 500bp with inward orientation

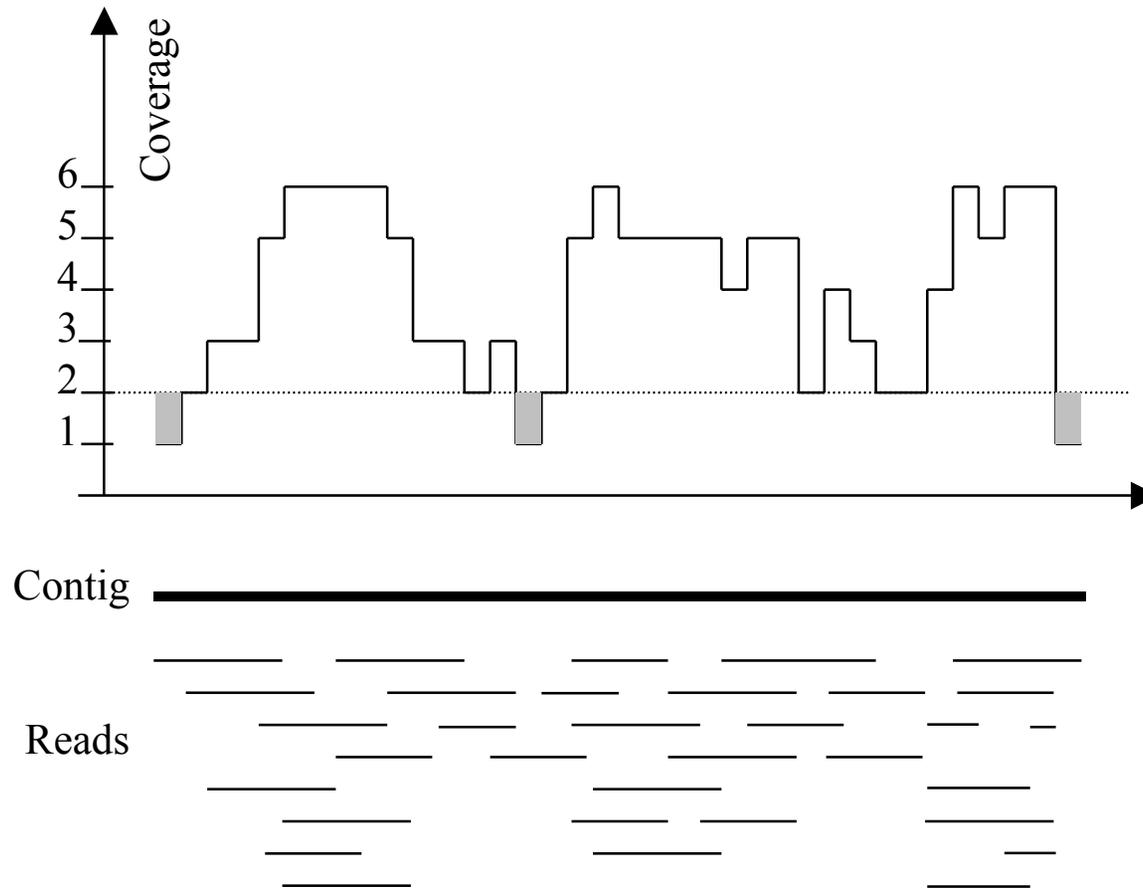


## ***Mate-pair sequencing***

- Circularize long molecules (1-10kbp), shear into fragments, & sequence
- Mate failures create short paired-end reads



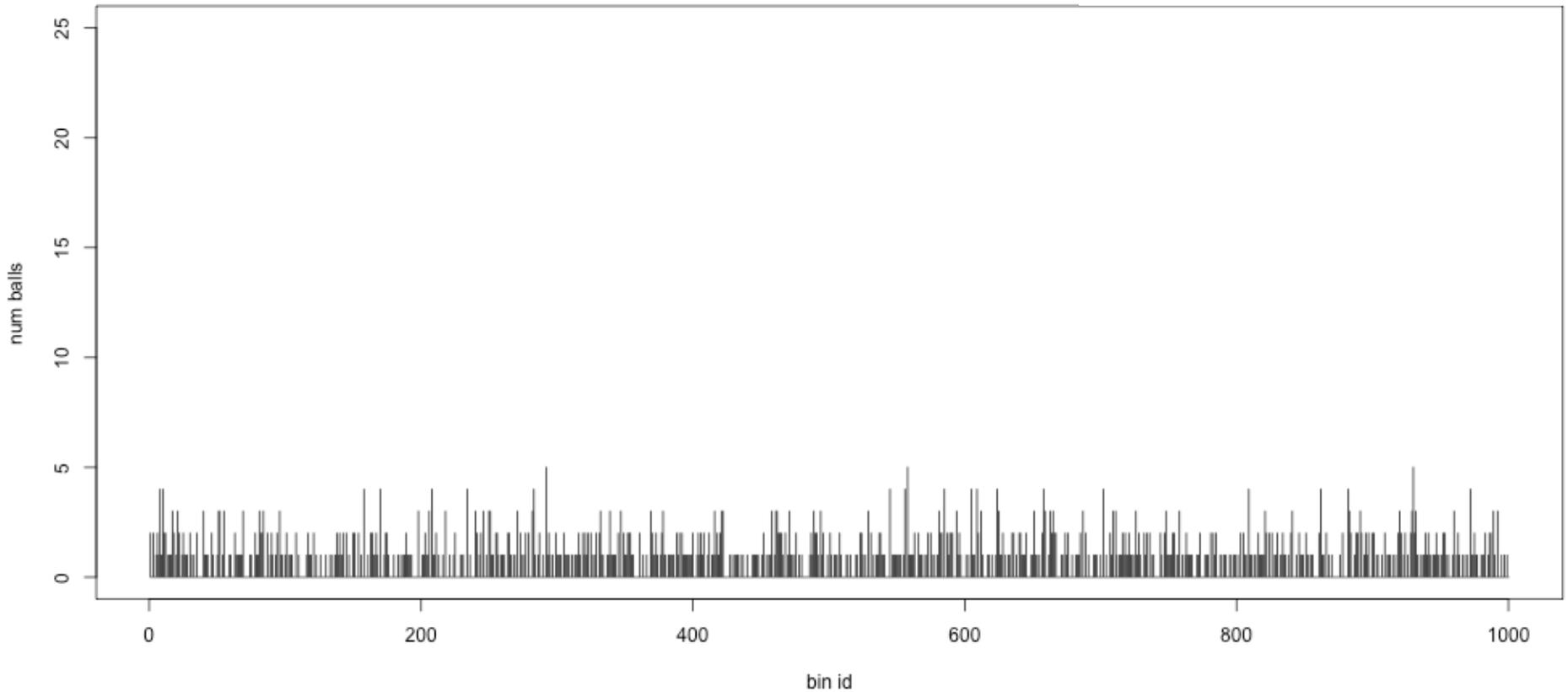
# Typical contig coverage



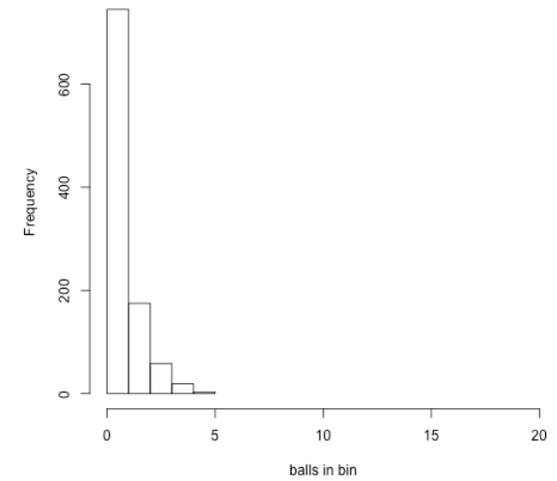
Imagine raindrops on a sidewalk

# Balls in Bins Ix

**Balls in Bins**  
Total balls: 1000

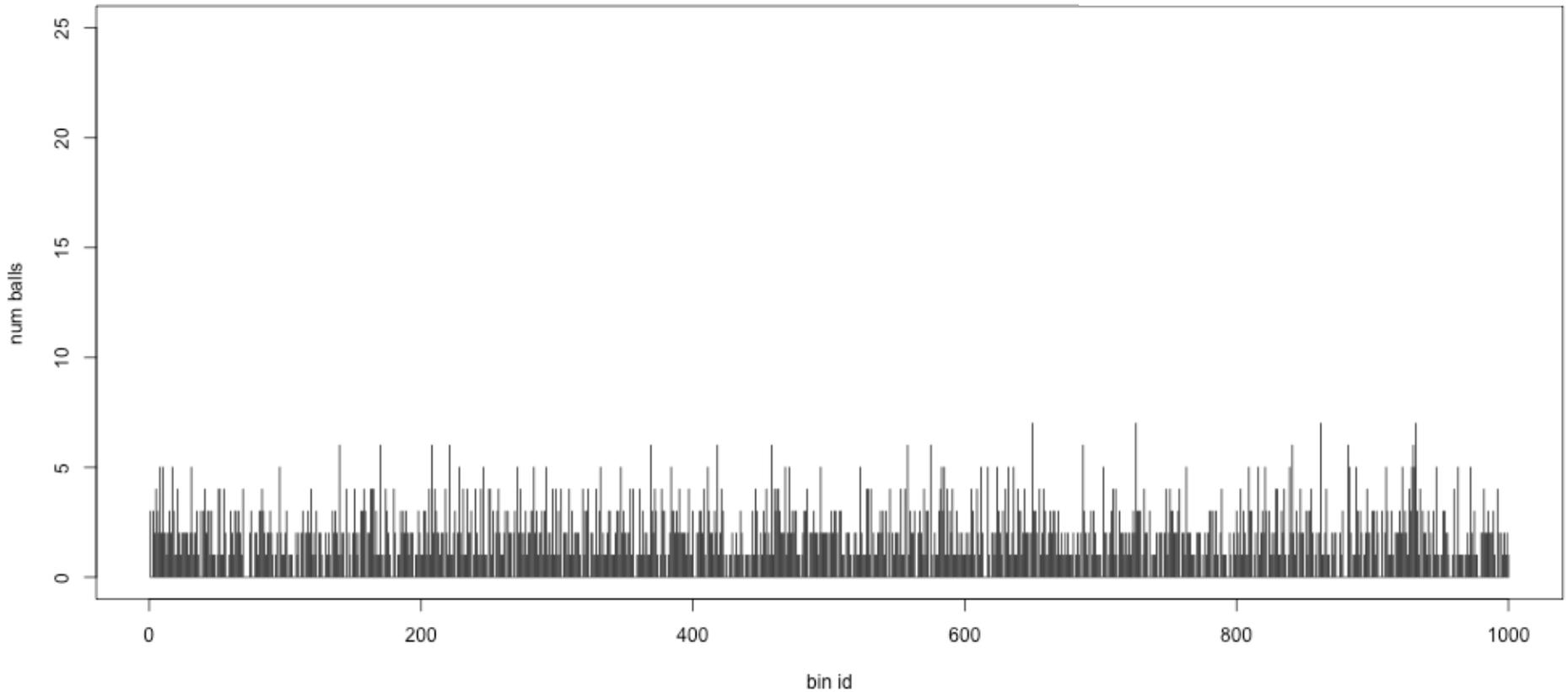


**Histogram of balls in each bin**  
Total balls: 1000

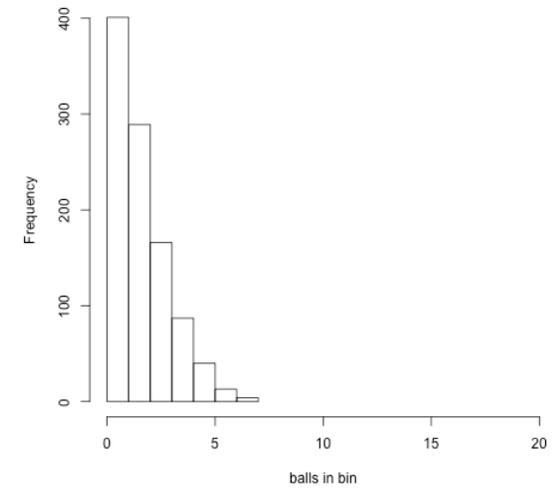


# Balls in Bins 2x

**Balls in Bins**  
Total balls: 2000

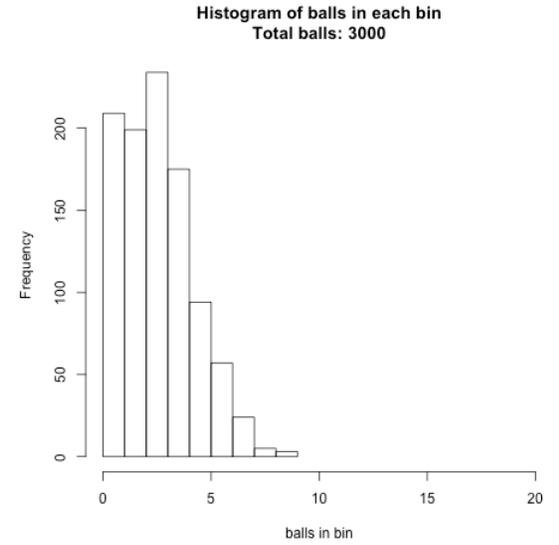
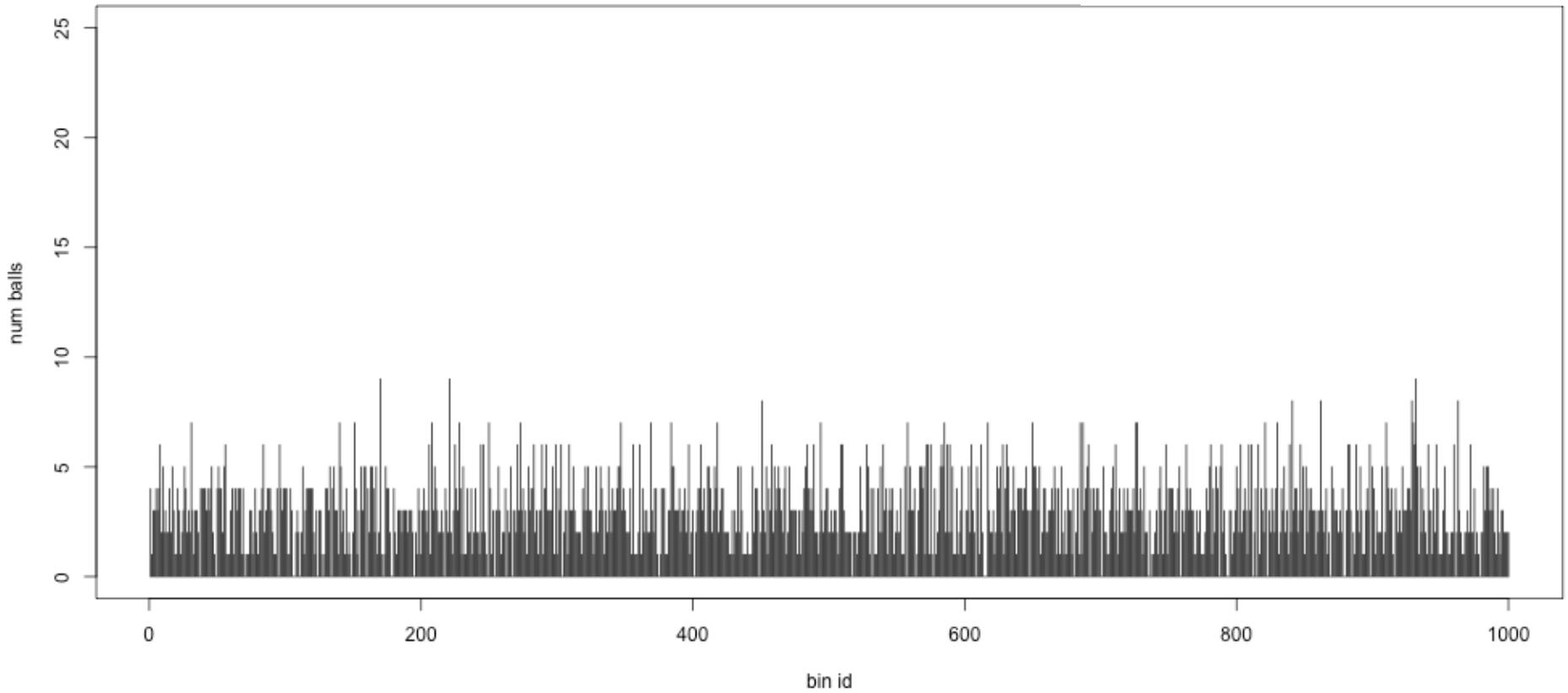


**Histogram of balls in each bin**  
Total balls: 2000



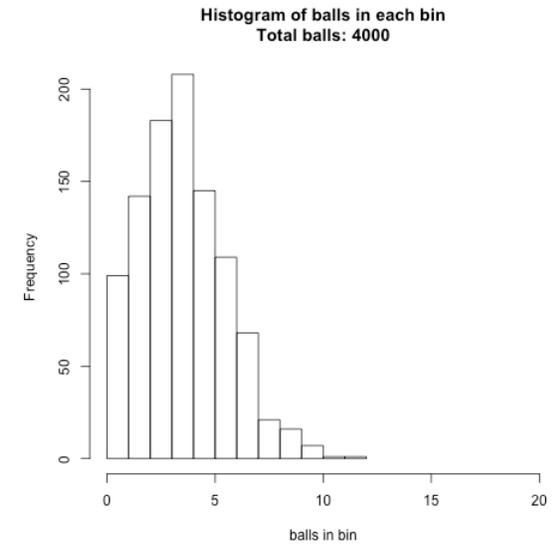
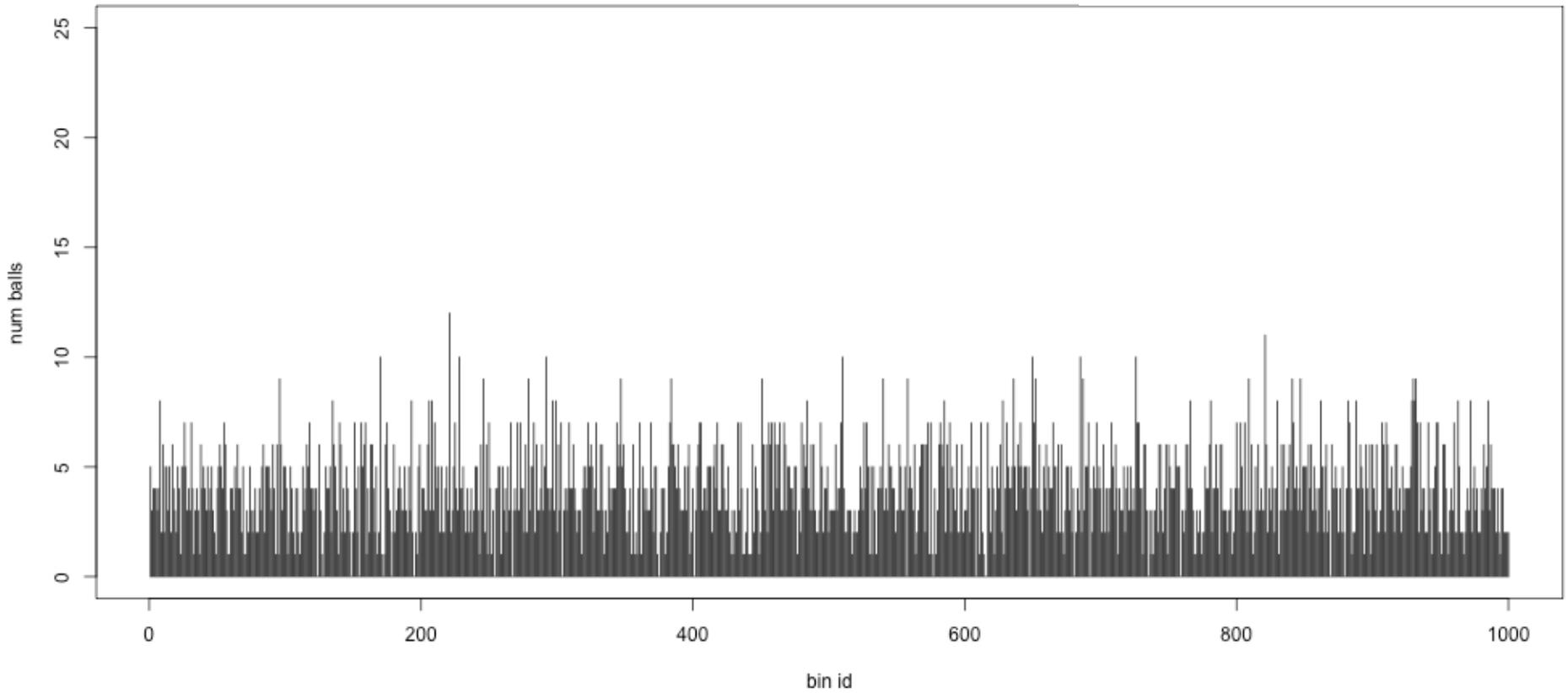
# Balls in Bins 3x

**Balls in Bins**  
Total balls: 3000



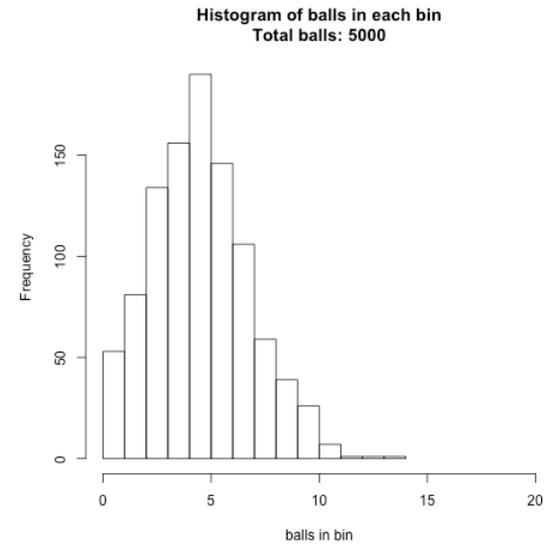
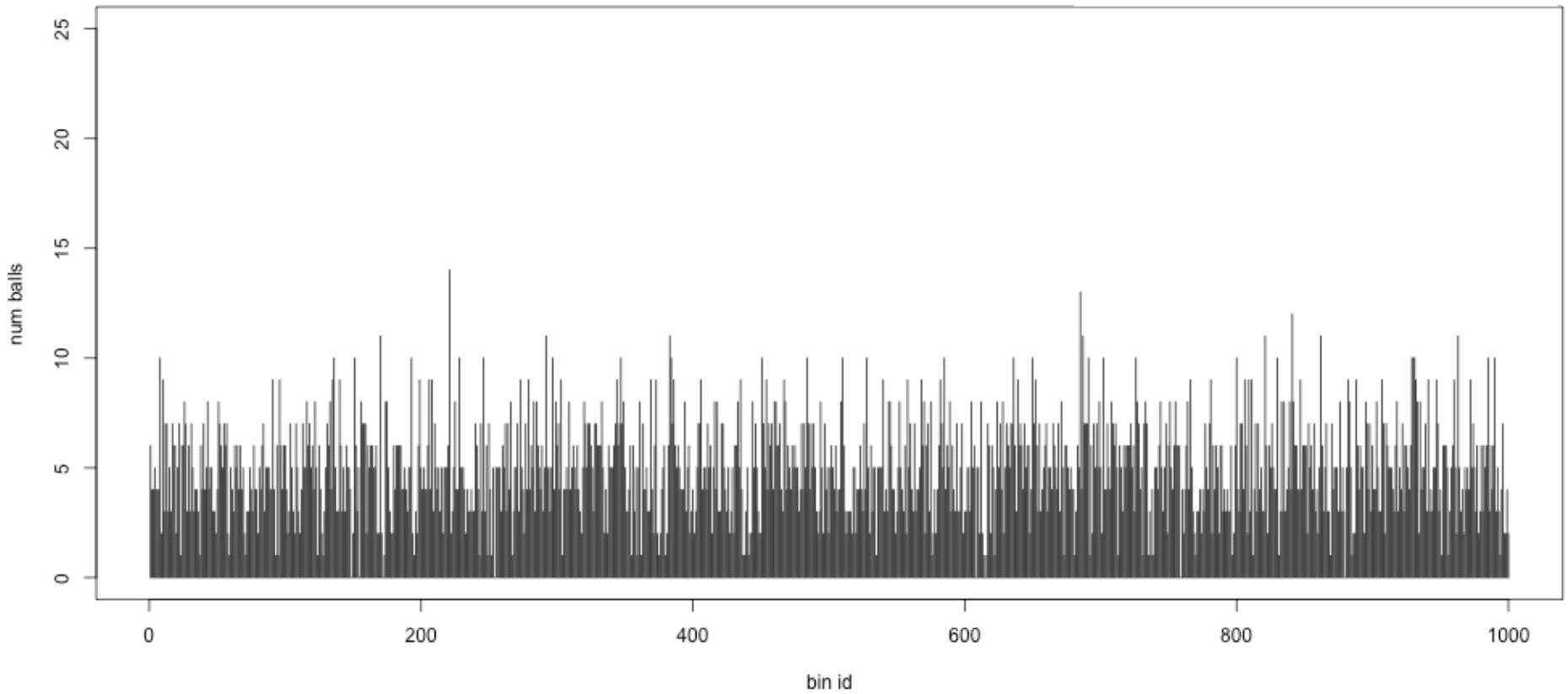
# Balls in Bins 4x

**Balls in Bins**  
Total balls: 4000



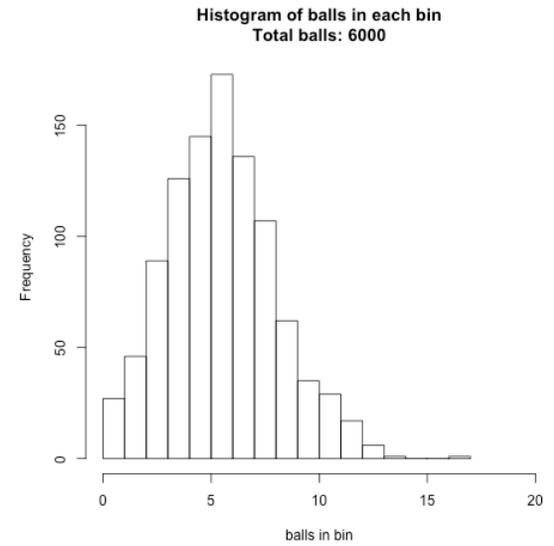
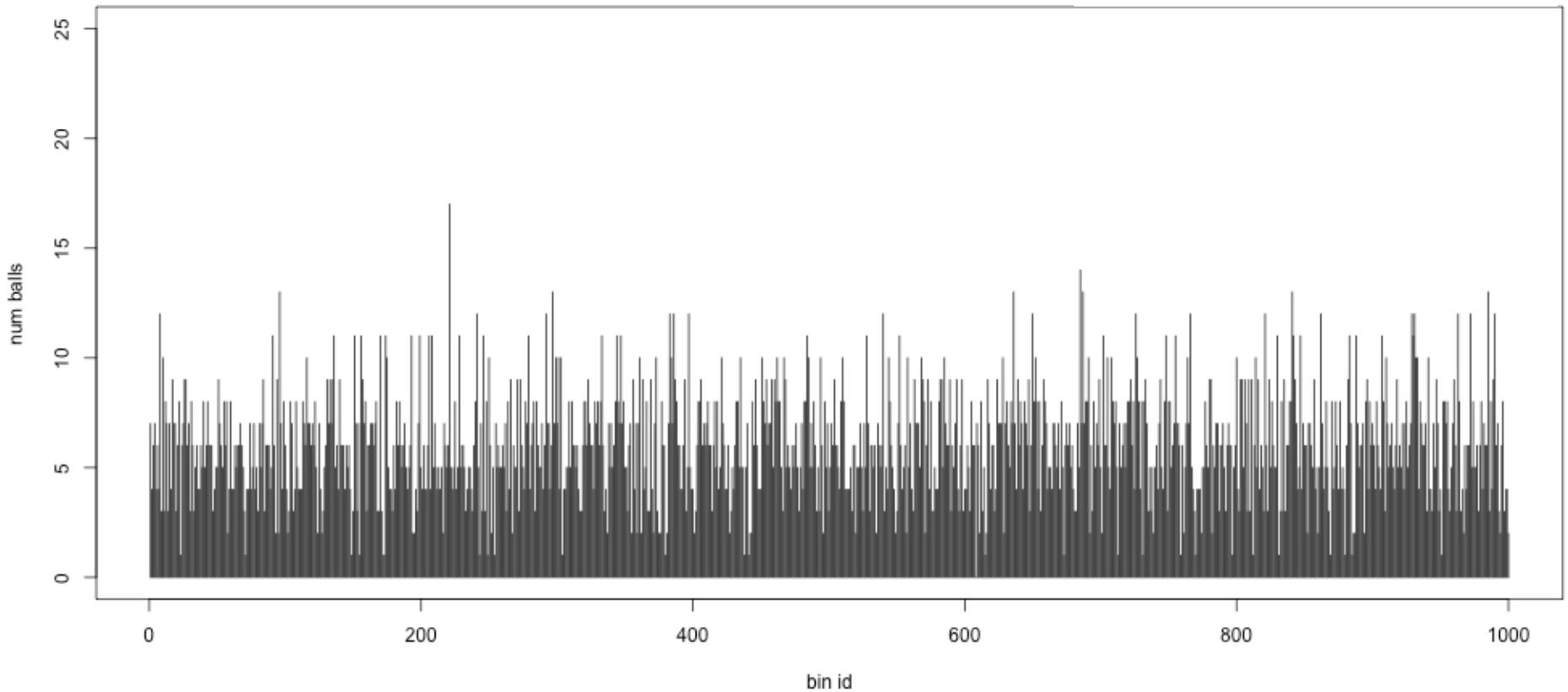
# Balls in Bins 5x

**Balls in Bins**  
Total balls: 5000



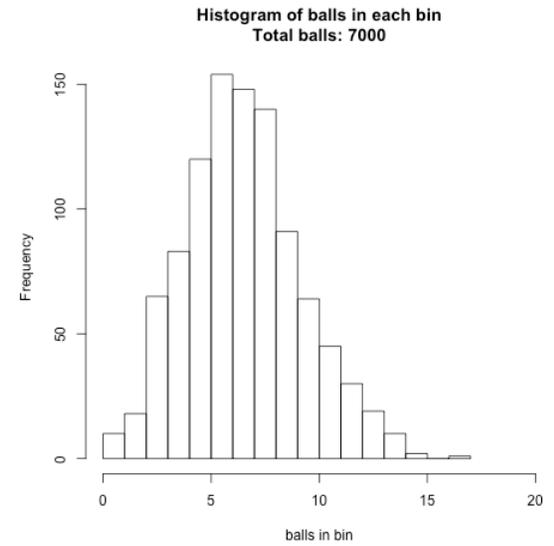
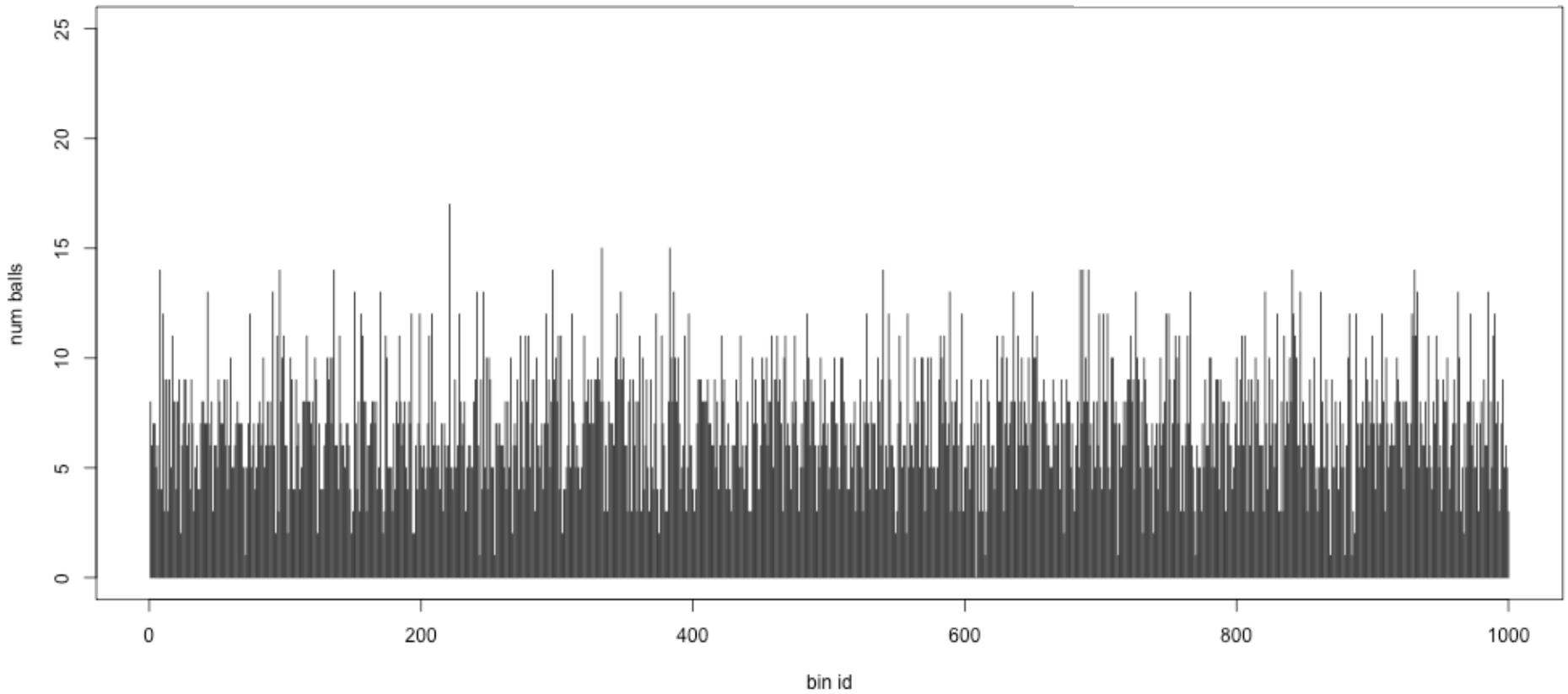
# Balls in Bins 6x

**Balls in Bins**  
Total balls: 6000



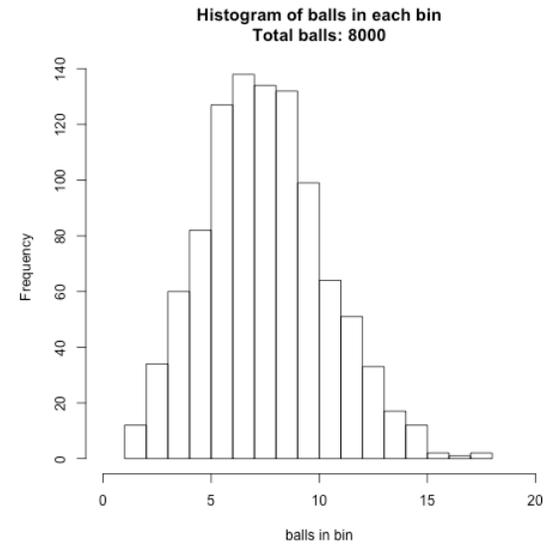
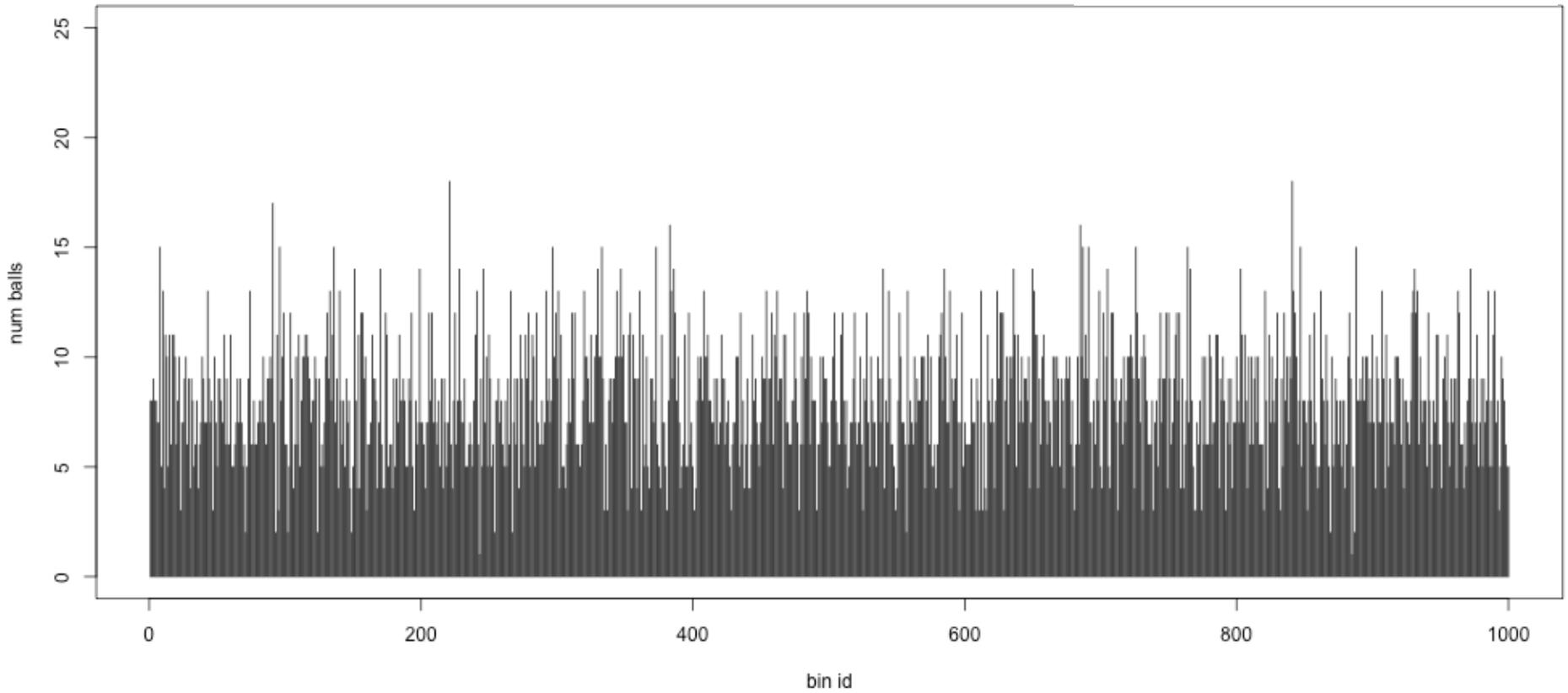
# Balls in Bins 7x

**Balls in Bins**  
Total balls: 7000



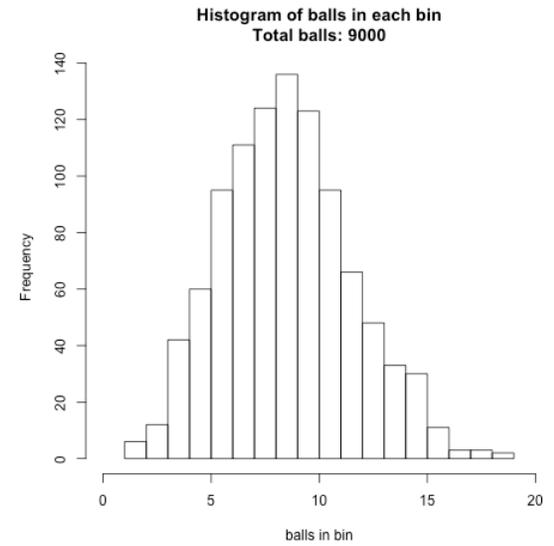
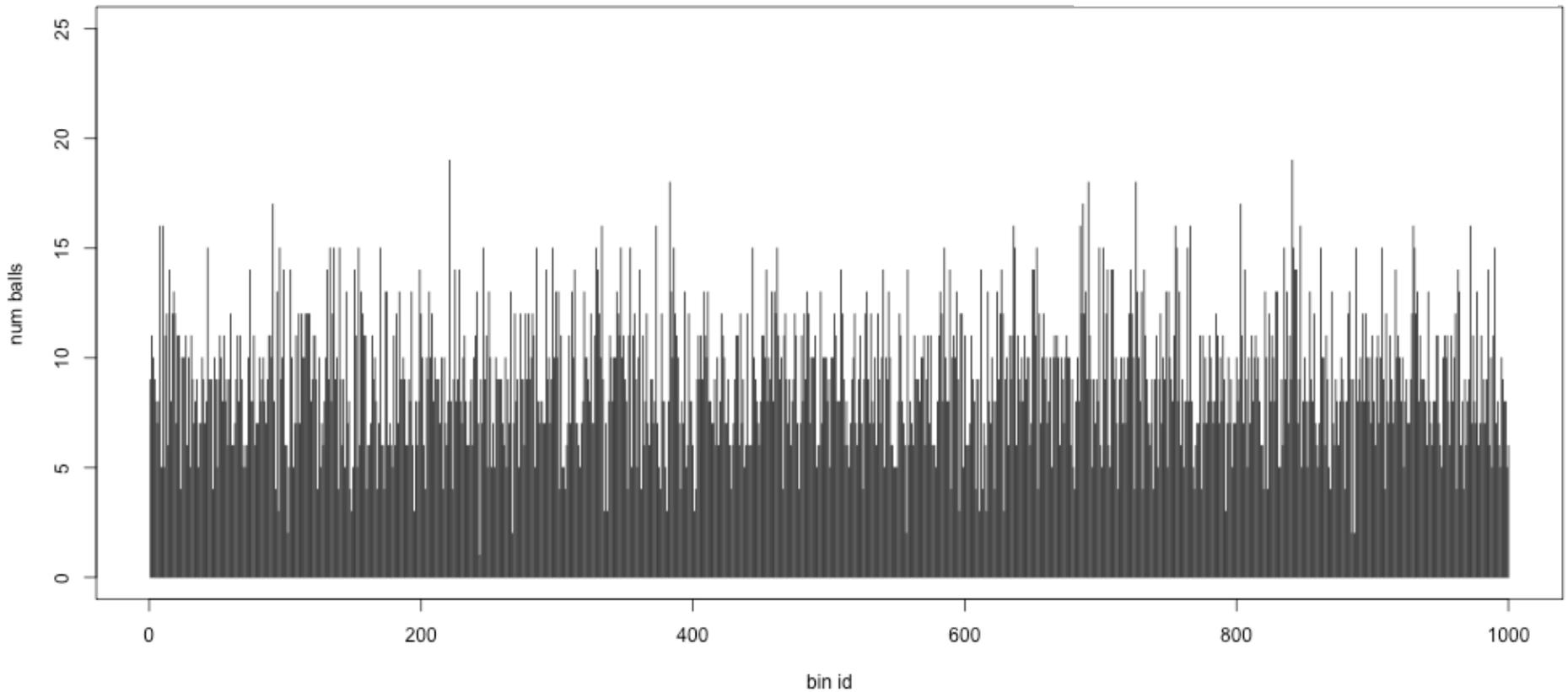
# Balls in Bins 8x

**Balls in Bins**  
Total balls: 8000



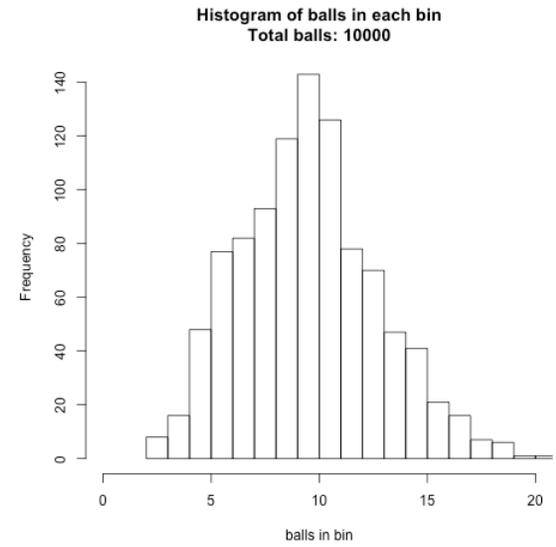
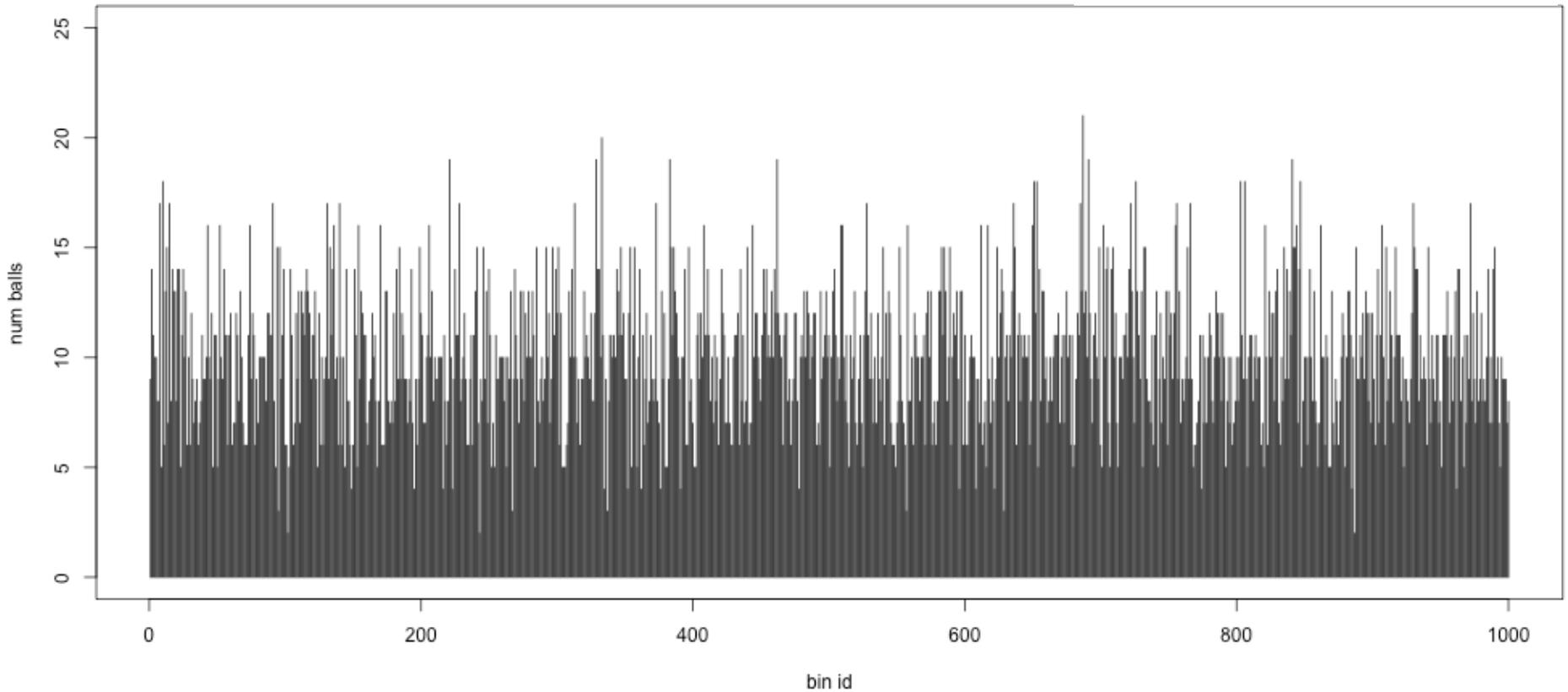
# Balls in Bins 9x

**Balls in Bins**  
Total balls: 9000

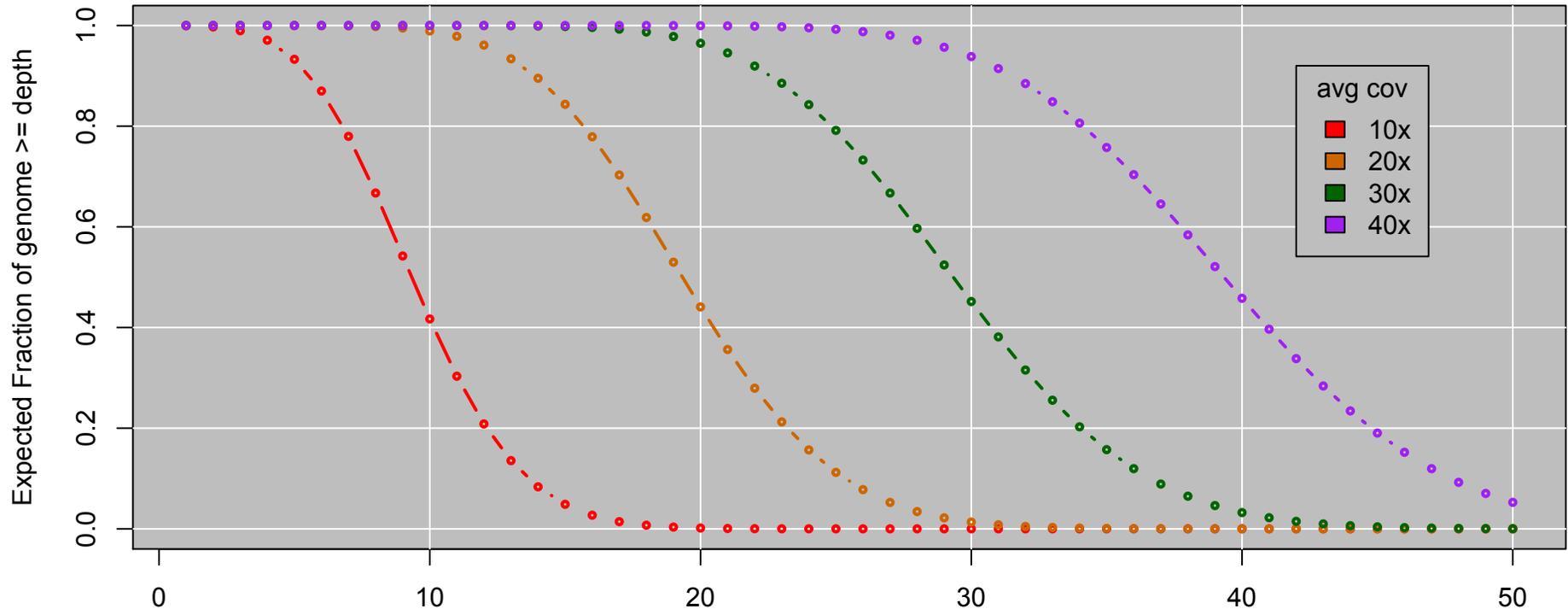


# Balls in Bins 10x

**Balls in Bins**  
Total balls: 10000



# Genome Coverage Distribution

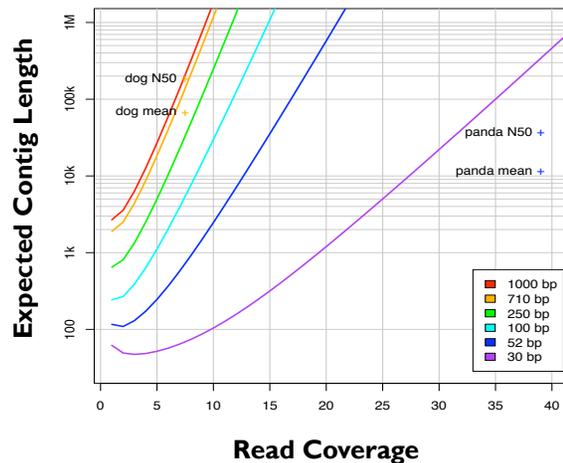


Expect Poisson distribution on depth  
Standard Deviation =  $\sqrt{\text{cov}}$

This is the mathematically model  $\Rightarrow$  reality may be much worse  
Double your coverage for diploid genomes

# Ingredients for a good assembly

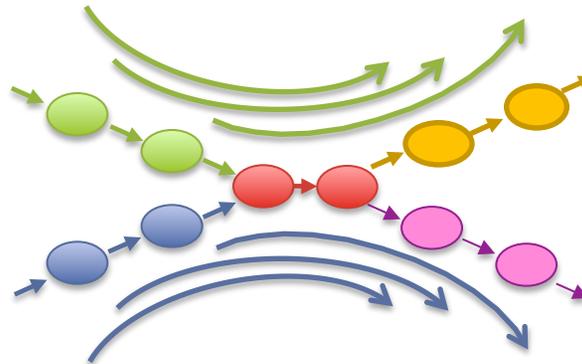
## Coverage



### High coverage is required

- Oversample the genome to ensure every base is sequenced with long overlaps between reads
- Biased coverage will also fragment assembly

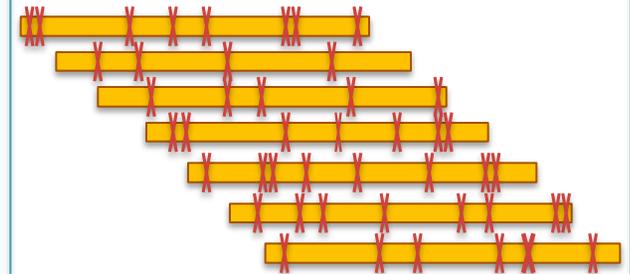
## Read Length



### Reads & mates must be longer than the repeats

- Short reads will have **false overlaps** forming hairball assembly graphs
- With long enough reads, assemble entire chromosomes into contigs

## Quality



### Errors obscure overlaps

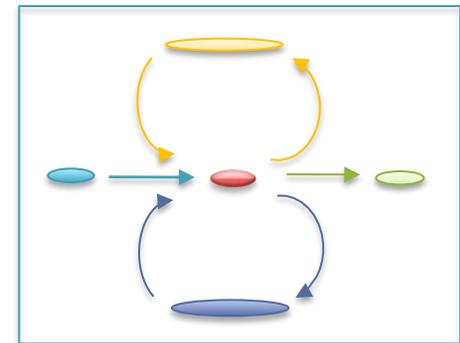
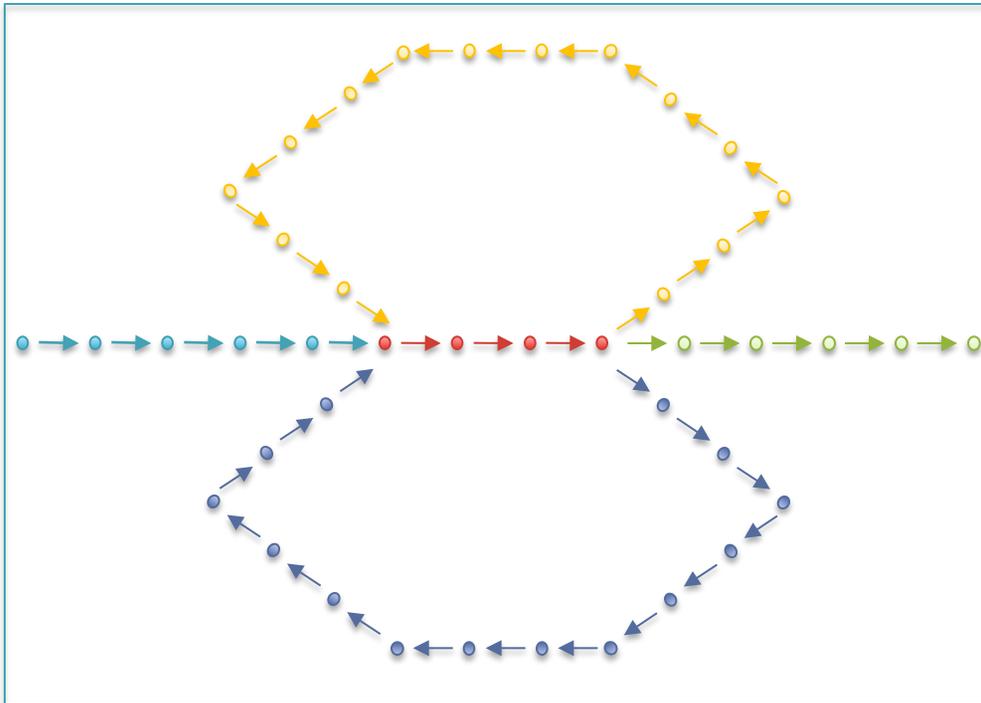
- Reads are assembled by finding kmers shared in pair of reads
- High error rate requires very short seeds, increasing complexity and forming assembly hairballs

## Current challenges in *de novo* plant genome sequencing and assembly

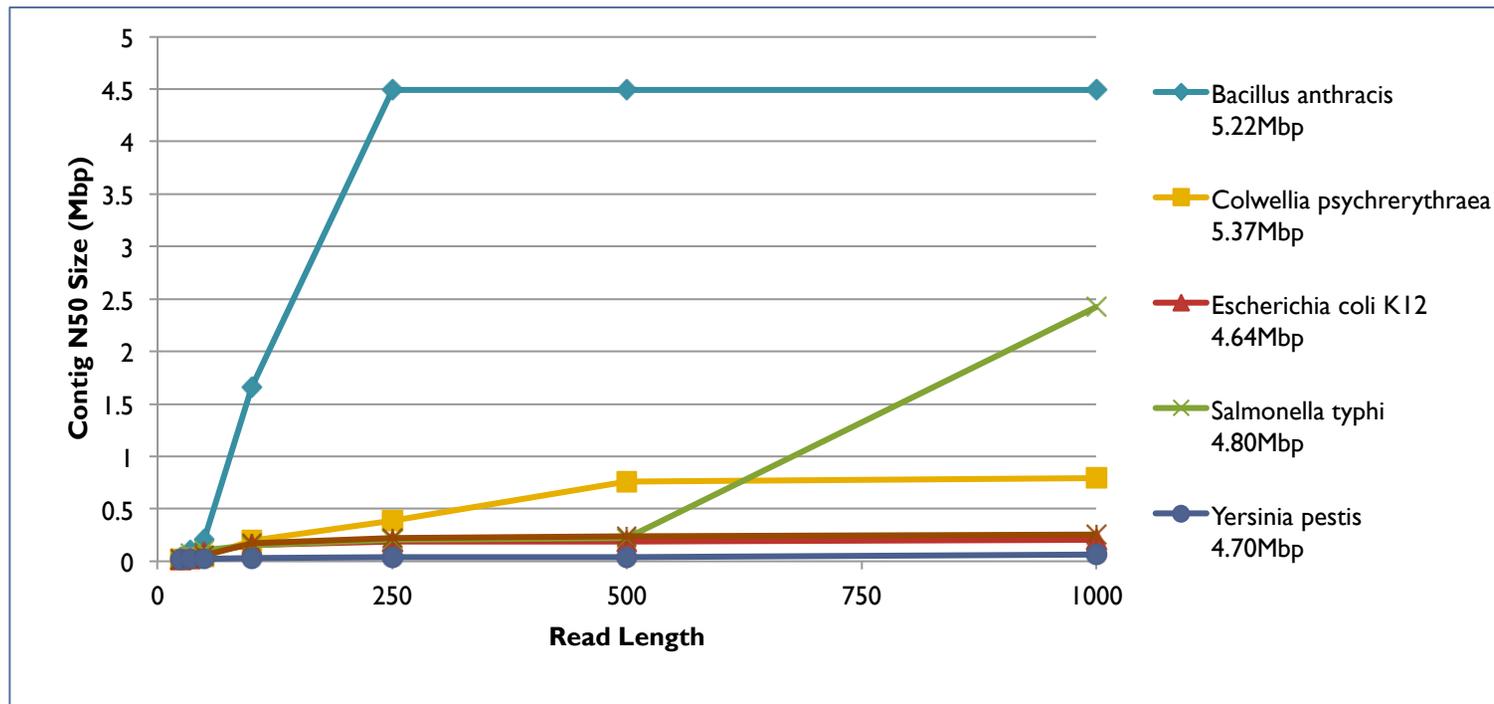
Schatz MC, Witkowski, McCombie, VWR (2012) *Genome Biology*.

# Initial Contigs

- After simplification and correction, compress graph down to its non-branching initial contigs
  - Aka “unitigs”, “unipaths”



# Repeats and Read Length



- Explore the relationship between read length and contig N50 size
  - Idealized assembly of read lengths: 25, 35, 50, 100, 250, 500, 1000
  - Contig/Read length relationship depends on specific repeat composition

## Assembly Complexity of Prokaryotic Genomes using Short Reads.

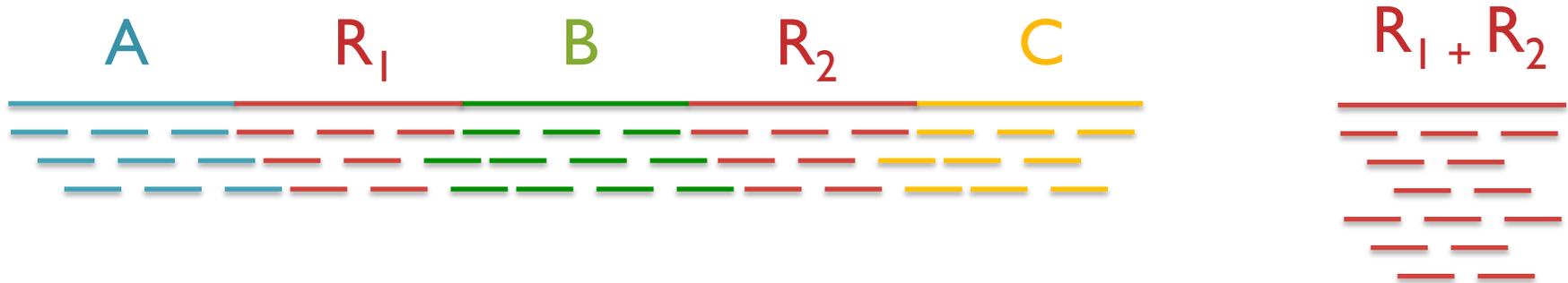
Kingsford C, Schatz MC, Pop M (2010) *BMC Bioinformatics*. 11:21.

# Repetitive regions

- Over 50% of the human genome is repetitive

| Repeat Type                                 | Definition / Example  | Prevalence |
|---|---|------------|
| Low-complexity DNA / Microsatellites        | $(b_1b_2\dots b_k)^N$ where $1 \leq k \leq 6$<br>CACACACACACACACACA | 2%         |
| SINEs (Short Interspersed Nuclear Elements) | <i>Alu</i> sequence (~280 bp)<br>Mariner elements (~80 bp)          | 13%        |
| LINEs (Long Interspersed Nuclear Elements)  | ~500 – 5,000 bp   | 21%        |
| LTR (long terminal repeat) retrotransposons | Ty1-copia, Ty3-gypsy, Pao-BEL<br>(~100 – 5,000 bp)                  | 8%         |
| Other DNA transposons                       |   | 3%         |
| Gene families & segmental duplications      |   | 4%         |

# Repeats and Coverage Statistics



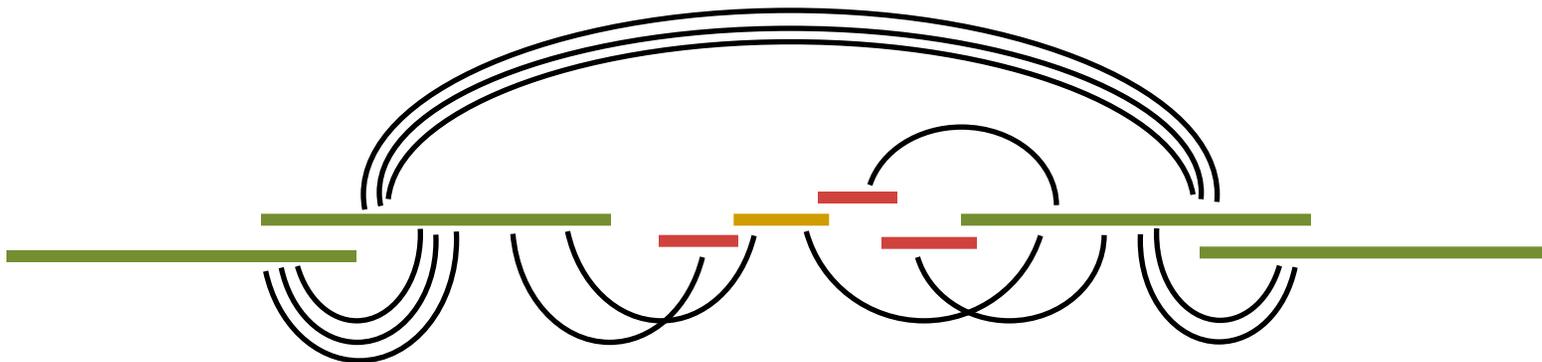
- If  $n$  reads are a uniform random sample of the genome of length  $G$ , we expect  $k = n \Delta / G$  reads to start in a region of length  $\Delta$ .
  - If we see many more reads than  $k$  (if the arrival rate is  $> \lambda$ ), it is likely to be a collapsed repeat
  - Requires an accurate genome size estimate

$$\Pr(X - \text{copy}) = \binom{n}{k} \left( \frac{\Delta n}{G} \right)^k \left( \frac{G - \Delta n}{G} \right)^{n-k}$$

$$A(\Delta, k) = \ln \left( \frac{\Pr(1 - \text{copy})}{\Pr(2 - \text{copy})} \right) = \ln \left( \frac{\frac{(\Delta n / G)^k e^{-\frac{\Delta n}{G}}}{k!}}{\frac{(2\Delta n / G)^k e^{-\frac{2\Delta n}{G}}}{k!}} \right) = \frac{n\Delta}{G} - k \ln 2$$

# Scaffolding

- Initial contigs (*aka* unipaths, unitigs) terminate at
  - *Coverage gaps*: especially extreme GC regions
  - *Conflicts*: sequencing errors, repeat boundaries
- Iteratively resolve longest, ‘most unique’ contigs
  - Both overlap graph and de Bruijn assemblers initially collapse repeats into single copies
  - Uniqueness measured by a statistical test on coverage

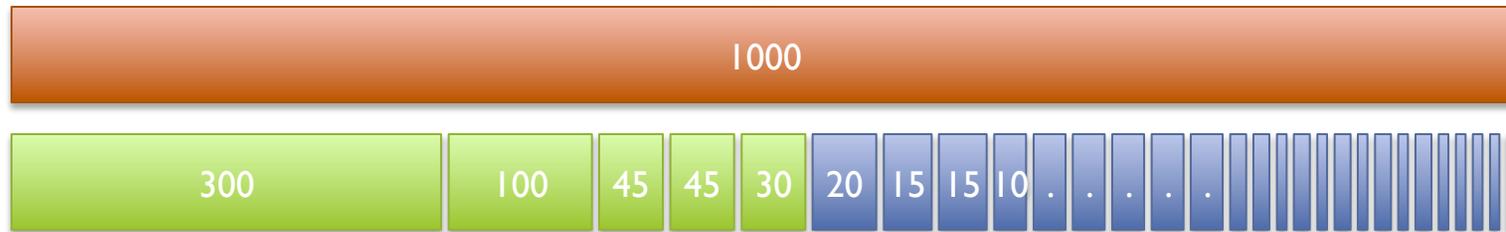


# N50 size

Def: 50% of the genome is in contigs larger than N50

Example: 1 Mbp genome

50%



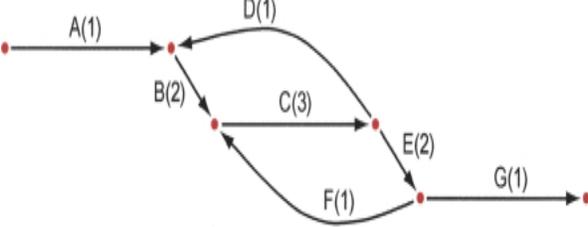
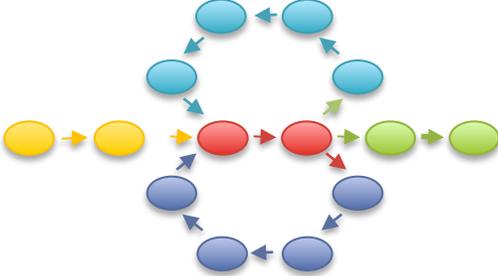
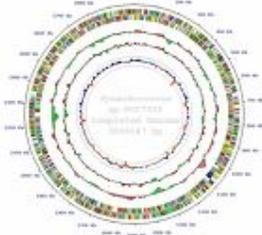
N50 size = 30 kbp

$(300k + 100k + 45k + 45k + 30k = 520k \geq 500kbp)$

Note:

N50 values are only meaningful to compare when base genome size is the same in all cases

# Assembly Algorithms

| ALLPATHS-LG   | SOAPdenovo  | Celera Assembler  |
|---|---|---|
|    |                                 |  |
| <p>Broad's assembler<br/>(Gnerre et al. 2011)</p>   | <p>BGI's assembler<br/>(Li et al. 2010)</p>   | <p>JCVI's assembler<br/>(Miller et al. 2008)</p>                                    |
| <p>De bruijn graph<br/>Short + PacBio (patching)</p>  | <p>De bruijn graph<br/>Short reads</p>  | <p>Overlap graph<br/>Medium + Long reads</p>  |
| <p>Easy to run if you have<br/>compatible libraries</p>   | <p>Most flexible, but requires a<br/>lot of tuning</p>  | <p>Supports Illumina/454/PacBio<br/>Hybrid assemblies</p>                           |
| <p><a href="http://www.broadinstitute.org/software/allpaths-lg/blog/">http://www.broadinstitute.org/<br/>software/allpaths-lg/blog/</a></p> | <p><a href="http://soap.genomics.org.cn/soapdenovo.html">http://soap.genomics.org.cn/<br/>soapdenovo.html</a></p> | <p><a href="http://wgs-assembler.sf.net">http://wgs-assembler.sf.net</a></p>        |

# PacBio Error Correction & Assembly

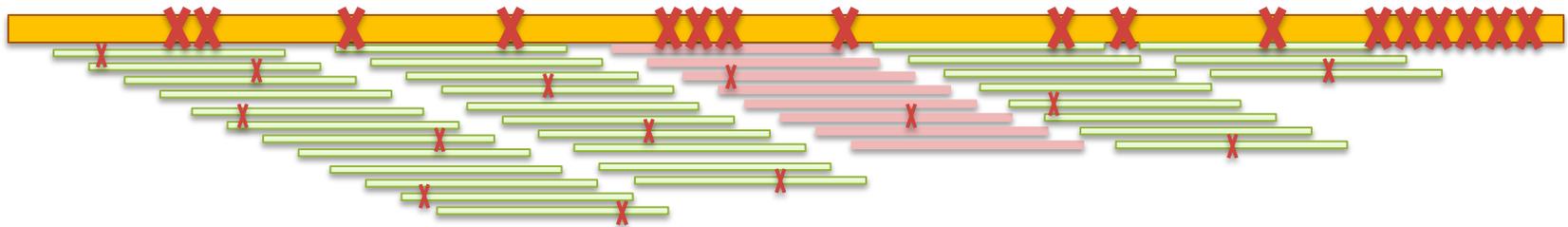
<http://wgs-assembler.sf.net>



## I. Correction Pipeline

1. Map short reads (SR) to long reads (LR)
2. Trim LR at coverage gaps
3. Compute consensus for each LR

## 2. Error corrected reads can be easily assembled, aligned



**Hybrid error correction and de novo assembly of single-molecule sequencing reads.**

Koren, S, Schatz, MC, Walenz, BP, Martin, J, Howard, J, Ganapathy, G, Wang, Z, Rasko, DA, McCombie, VWR, Jarvis, ED, Phillippy, AM. (2012) *Nature Biotechnology*. doi:10.1038/nbt.2280

# Assembly of Heterozygous Genomes

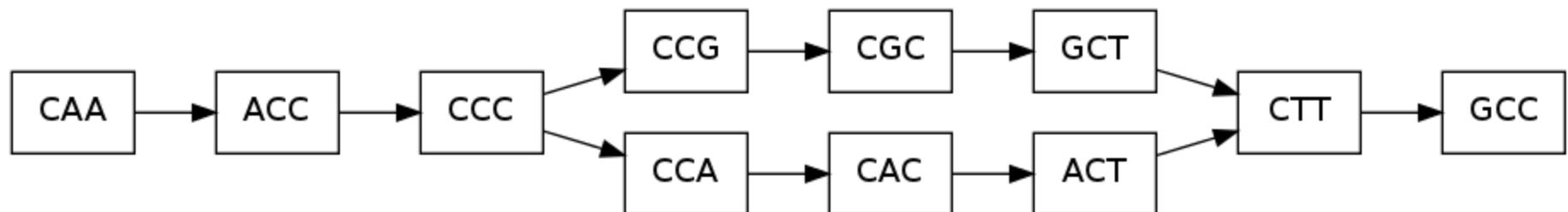
E. Biggers, M. Schatz

Genome assemblers developed to assembly genomes with low rates of heterozygosity

- 0-.1% (similar to human)

Assembly becomes more complicated with higher rates

Preprocess the reads to “smooth” the heterozygosity, assemble, and then restore variants

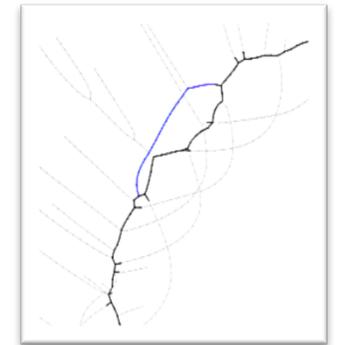


# Scalpel: Haplotype Microassembly

G. Narzisi, D. Levy, I. Iossifov, J. Kendall, M. Wigler, M. Schatz



- Use assembly techniques to identify complex variations from short reads
  - Improved power to find indels
  - Trace candidate haplotypes sequences as paths through assembly graphs



Ref: ...TCAGAACAGCTGGATGAGATCTTAGCCAACTACCAGGAGATTGTCTTTGCCCGGA...

Father: ...TCAGAACAGCTGGATGAGATCTTAGCCAACTACCAGGAGATTGTCTTTGCCCGGA...

Mother: ...TCAGAACAGCTGGATGAGATCTTAGCCAACTACCAGGAGATTGTCTTTGCCCGGA...

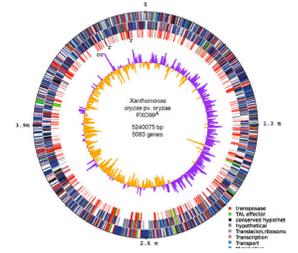
Sib: ...TCAGAACAGCTGGATGAGATCTTAGCCAACTACCAGGAGATTGTCTTTGCCCGGA...

Aut(1): ...TCAGAACAGCTGGATGAGATCTTAGCCAACTACCAGGAGATTGTCTTTGCCCGGA...

Aut(2): ...TCAGAACAGCTGGATGAGATCTTACC-----CCGGGAGATTGTCTTTGCCCGGA...

6bp heterozygous indel at chr13:25280526 ATP12A

# Assembly Summary



Graphs are ubiquitous in the world

- Pairwise searching is easy, finding features is hard

Assembly quality depends on

1. **Coverage**: low coverage is mathematically hopeless
2. **Repeat composition**: high repeat content is challenging
3. **Read length**: longer reads help resolve repeats
4. **Error rate**: errors reduce coverage, obscure true overlaps

Assembly is a hierarchical, starting from individual reads, build high confidence contigs/unitigs, incorporate the mates to build scaffolds

- Extensive error correction is the key to getting the best assembly possible from a given data set

# Genomics Challenges



**The foundations of genomics will continue to be *observation, experimentation, and interpretation***

- Technology will continue to push the frontier
- Measurements will be made *digitally* over large populations, at extremely high resolution, and for diverse applications

***Rise in Quantitative and Computational Demands***

1. *Experimental design*: selection, collection & metadata
2. *Observation*: measurement, storage, transfer, computation
3. *Integration*: multiple samples, assays, analyses
4. *Discovery*: visualizing, interpreting, modeling

***Ultimately limited by the human capacity to execute extremely complex experiments and interpret results***

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Lippman Lab

Martienssen Lab

McCombie Lab

Ware Lab

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Cole Trapnell



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<http://schatzlab.cshl.edu/teaching/>  
[@mike\\_schatz](#)